

CS 4410

Automata, Computability, and  
Formal Language

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# Chapter 8

## Properties of Context-Free Languages

1. Two Pumping Lemmas
  - A Pumping Lemma for Context-Free Languages
  - A Pumping Lemma for Linear Language
2. Closure Properties and Decision Algorithms for Context-Free Languages
  - Closure of Context-Free Languages
  - Some Decidable Properties of Context-Free Languages

# Closure of Context-Free Languages

**Theorem 8.5:** Let  $L_1$  be a context-free language and  $L_2$  be a regular language. Then  $L_1 \cap L_2$  is context-free.

**Example 8.7:** The language  $L = \{a^n b^n : n \geq 0, n \neq 100\}$  is context-free.

# Closure of Context-Free Languages

**Example 8.8:** Show that the language

$$L = \{w \in \{a,b,c\}^* : n_a(w) = n_b(w) = n_c(w)\}$$

is not context-free.

# Elementary Questions about Context-Free Languages

- Given a context-free language  $L$  and an arbitrary string  $w$ , is there an algorithm to determine whether or not  $w$  is in  $L$ ?
- Given a context-free language  $L$ , is there an algorithm to determine if  $L$  is empty?
- Given a context-free language  $L$ , is there an algorithm to determine if  $L$  is infinite?
- Given two context-free grammars  $G_1$  and  $G_2$ , is there an algorithm to determine if  $L(G_1) = L(G_2)$ ?

# A Membership Algorithm for Context-Free Languages

- The combination of Theorems 5.2 and 6.5 confirms the existence of a membership algorithm for context-free languages
  - By Theorem 5.2, exhaustive parsing is guaranteed to give the correct result for any context-free grammar that contains neither  $\lambda$ -productions nor unit-productions
  - By Theorem 6.5, such a grammar can always be produced if the language does not include  $\lambda$
- Alternatively, a npda to accept the language can be constructed as established by Theorem 7.1

# Determining Whether a Context-Free Language is Empty

- **Theorem 8.6** confirms the existence of an algorithm to determine if a context-free language  $L(G)$  is empty
  - For simplicity, assume that  $\lambda$  is not in  $L(G)$
  - Apply the algorithm for removing useless symbols and productions
  - If the start symbol is found to be useless, then  $L(G)$  is empty; otherwise,  $L(G)$  contains at least one string

# Determining Whether a Context-Free Language is Infinite

- **Theorem 8.7** confirms the existence of an algorithm to determine if a context-free language  $L(G)$  is infinite
  - Apply the algorithms for removing  $\lambda$ -productions, unit-productions, and useless productions
  - If  $G$  has a variable  $A$  for which there is a derivation that allows  $A$  to produce a sentential form  $xAy$ , then  $L(G)$  is infinite. Otherwise,  $L(G)$  is finite



# Determining Whether Two Context-Free Languages are Equal

- Given two context-free grammars  $G_1$  and  $G_2$ , is there an algorithm to determine if  $L(G_1) = L(G_2)$ ?
- If the languages are finite, the answer can be found by performing a string-by-string comparison
- However, for general context-free languages, no algorithm exists to determine equality