

COOL

Cool Overview

- Classroom Object Oriented Language
- Designed to
 - Be implementable in a short time
 - Give a taste of implementation of modern
 - Abstraction
 - Static typing
 - Reuse (inheritance)
 - Memory management
- But many things are left out

A Simple Example

```
class Point {  
    x : Int ← 1;  
    y : Int ← 2;  
};
```

- Cool programs are sets of class definitions
 - class = a collection of attributes and methods
 - instances of a class are objects
- No global variables
- No separate notion of subroutines
 - Entry point is a special class **Main** with a special method **main**

Cool Objects

```
class Point {  
    x : Int ← 3;  
    y : Int; (* use default value *)  
};
```

- The expression “**new Point**” creates a new instance (i.e. object) of class **Point**
- An object can be thought of as a record with a slot for each attribute

x	y
3	0

Methods

- A class defines methods for manipulating the attributes

```
class Point {  
  x : Int ← 0;  
  y : Int ← 0;  
  movePoint(newx : Int, newy : Int): Point {  
    { x ← newx;  
      y ← newy;  
      self;  
    } -- close block expression  
  }; -- close method  
}; -- close class
```

- Methods can refer to the current object using `self`

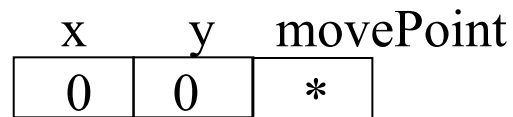
Information Hiding in Cool

- Methods are global
- Attributes are local to a class
 - They can only be accessed by the class's methods
- Example:

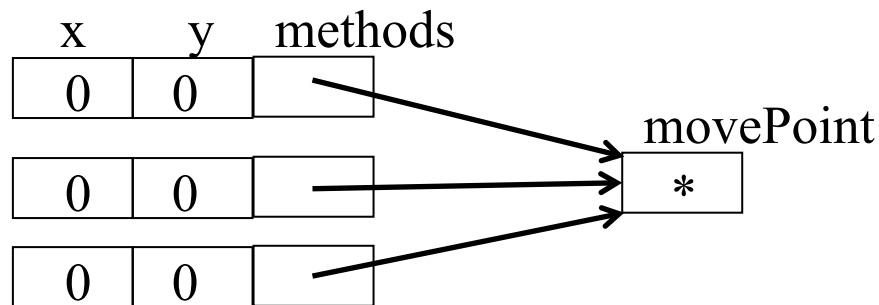
```
class Point {  
    . . .  
    x () : Int { x };  
    setx (newx : Int) : Int { x ← newx };  
};
```

Methods

- Each object knows how to access the code of a method
- As if the object contains a slot pointing to the code



- In reality implementations save space by sharing these pointers among instances of the same class



Inheritance

- We can extend functionality of points to colored points using subclassing

```
class ColorPoint inherits Point {  
  color : String ← "red";  
  movePoint(newx : Int, newy : Int) : Point {  
    { color ← "green";  
      x ← newx; y ← newy;  
      self;  
    }  
  };  
};
```

x	y	color	movePoint
0	0	red	*

Cool Types

- Every class is a type
- Every* class inherits from exactly one other class
 - Forms a tree of classes (*class hierarchy*)
- Types of all variables must be declared
 - Compiler infers types for expressions

Base Classes

- **Object** root of the class hierarchy
- **Int** integers
- **Bool** boolean values: **true**, **false**
- **String** character strings
- **IO** input/output support

Cool Type Checking

```
x : A;  
x ← new B;
```

- Is *well-typed* if **A** is an ancestor of **B** in the class hierarchy
 - Anywhere an instance of **A** is expected an instance of **B** can be used
- “Well-typed” = satisfies language’s type-checking rules
- Type safety:
 - A well-typed program cannot result in runtime type errors

Method Invocation and Inheritance

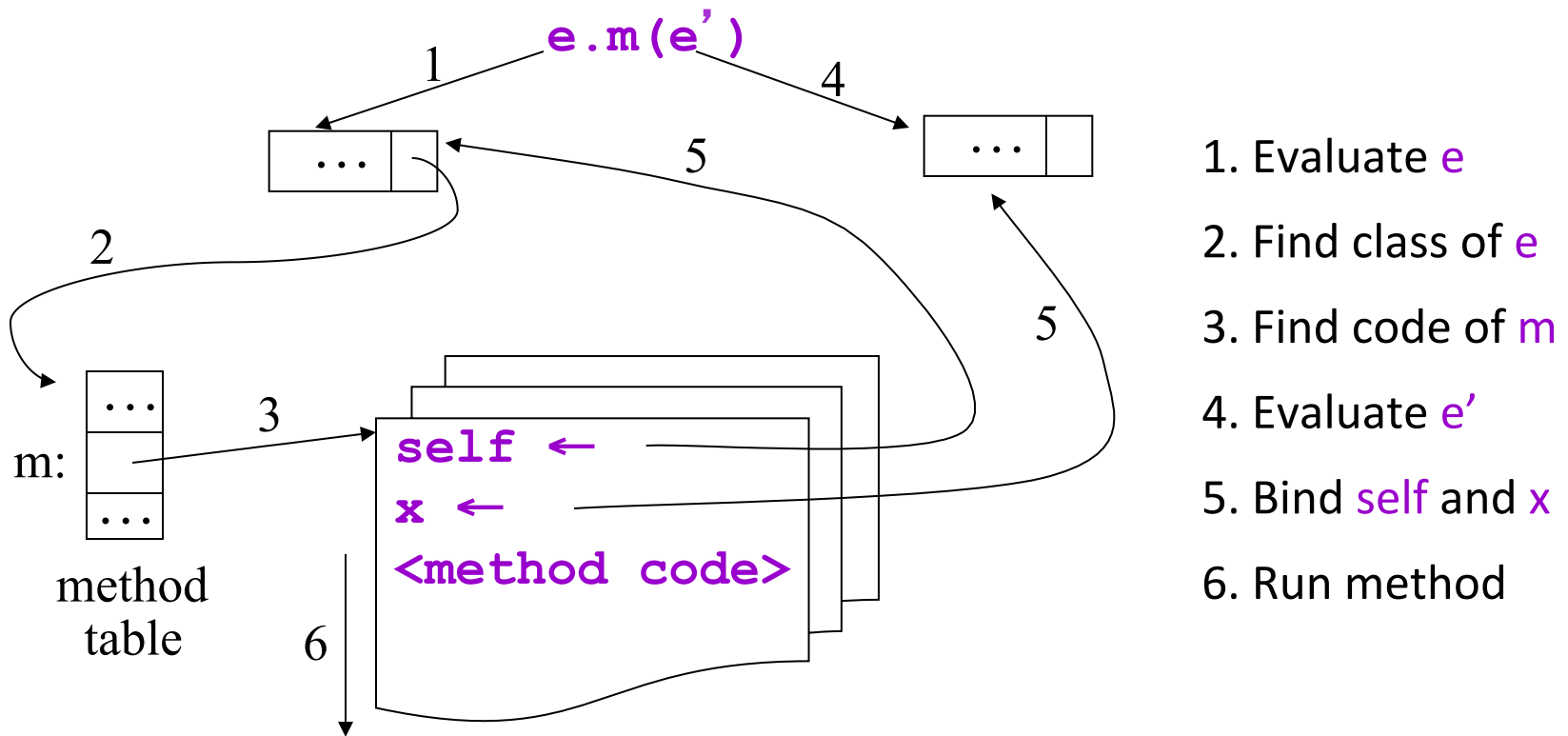
- Methods are invoked by *dispatching* them to the target object
- Understanding dispatch in the presence of inheritance is a subtle aspect of OO languages

```
p : Point;  
p ← new ColorPoint;  
p.movePoint(1,2);
```

- `p` has static type `Point`
- `p` has dynamic type `ColorPoint`
- `p.movePoint` must invoke the `ColorPoint` version

Method Invocation

- Example: invoke one-argument method m



Other Expressions

- Expression language
 - every expression has a type and a value
 - Loops: `while E loop E pool`
 - Conditionals `if E then E else E fi`
 - Case statement `case E of x : Type => E; ... esac`
 - Arithmetic, logical operations
 - Assignment `x ← E`
 - Primitive I/O `out_string(s), in_string(), ...`
- Missing features:
 - arrays, floating point operations, exceptions, ...

Other Expressions (Cont.)

- Blocks
 - { <expr>; ... <expr>; }
- Let
 - let <id1> : <type1> [<- <expr1>], ..., <idn> : <typen> [<- <exprn>] in <expr>

Cool Memory Management

- Memory is allocated every time `new` is invoked
- Memory is deallocated automatically when an object is not reachable anymore
 - Done by the garbage collector (GC)
 - Garbage collector is part of the Cool runtime