## This is the second lecture of Chapter 11



Measurement and Analysis (B)

## Quick review of last lecture

- Introduction
- Computer Performance Equations
- CPU time
- Amdahl's Law
- Quantitative Principles of Computer Design
- Amdahl's Law and Parallel Speedup
- The CPU Performance Equation
- Average CPI, IC, Clock Cycle Time, Clock Rate


### 11.3 Mathematical Preliminaries (1 of 17)

- Measures of system performance depend upon one's point of view.
- A computer user is most often concerned with response time: How long does it take the system to carry out a task?
- System administrators are usually more concerned with throughput: How many concurrent tasks can the system handle before response time is adversely affected?
- These two ideas are related: If a system carries out a task in $k$ seconds, then its throughput is $1 / k$ of these tasks per second.


### 11.3 Mathematical Preliminaries (2 of 17)

- In comparing the performance of two systems, we measure the time that it takes for each system to do the same amount of work.
- Specifically, if System A and System B run the same program, System A is $n$ times as fast as System B if:

$$
\frac{\text { running time on system } \mathrm{B}}{\text { running time on system } \mathrm{A}}=n
$$

- System A is x\% faster than System B if:

$$
\left[\frac{\text { running time on system } B}{\text { running time on system } A}-1\right] \times 100 \%=x \%
$$

### 11.3 Mathematical Preliminaries (3 of 17)

- Suppose we have two racecars that have just completed a 10 mile race. Car A finished in 3 minutes, and Car B finished in 4 minutes. Using our formulas, Car A is 1.33 times as fast as Car B, and Car A is also $33 \%$ faster than Car B:

$$
\frac{\text { time for Car B to travel } 10 \text { miles }}{\text { time for Car A to travel } 10 \text { miles }}=\frac{4}{3}=1.33 \text {. }
$$

$\left[\begin{array}{lll}\frac{\text { running time on system } B}{\text { running time on system } A} & -1\end{array}\right] \times 100 \%=\left(\frac{4}{3}-1\right) \times 100=33 \%$

### 11.3 Mathematical Preliminaries (4 of 17)

- When we are evaluating system performance we are most interested in its expected performance under a given workload.
- We use statistical tools that are measures of central tendency.
- The one with which everyone is most familiar is the arithmetic mean (or average), given by:

$$
\sum_{i=1}^{n} \mathbf{X}_{i}
$$

### 11.3 Mathematical Preliminaries (5 of 17)

- The arithmetic mean can be misleading if the data are skewed or scattered.
- Consider the execution times given in the table below. The performance differences are hidden by the simple average.

| Program | System A <br> Execution <br> Time | System B <br> Execution <br> Time | System C <br> Execution <br> Time |
| :---: | :---: | :---: | :---: |
| $v$ | 50 | 100 | 500 |
| $w$ | 200 | 400 | 600 |
| $x$ | 250 | 500 | 500 |
| $y$ | 400 | 800 | 800 |
| $z$ | 5,000 | 4,100 | 3,500 |
| Average | 1,180 | 1,180 | 1,180 |

### 11.3 Mathematical Preliminaries (6 of 17)

- If execution frequencies (expected workloads) are known, a weighted average can be revealing.
- The weighted average for System A is:
- $50 \times 0.5+200 \times 0.3+250 \times 0.1+400 \times 0.05+5000 \times$ $0.05=380$.

| Program | Execution <br> Frequency | System A <br> Execution <br> Time | System C <br> Execution <br> Time |
| :---: | :---: | :---: | :---: |
| $v$ | $50 \%$ | 50 | 500 |
| $w$ | $30 \%$ | 200 | 600 |
| $x$ | $10 \%$ | 250 | 500 |
| $y$ | $5 \%$ | 400 | 800 |
| $z$ | $5 \%$ | 5,000 | 3,500 |
| Weighted Average |  | 380 seconds | 695 seconds |

### 11.3 Mathematical Preliminaries

(7 of 17)

- However, workloads can change over time.
- A system optimized for one workload may perform poorly when the workload changes, as illustrated below.

| Program | Execution <br> Time | Execution <br> Frequency |
| :---: | :---: | :---: |
| $v$ | 50 | $25 \%$ |
| $w$ | 200 | $5 \%$ |
| $x$ | 250 | $10 \%$ |
| $y$ | 400 | $5 \%$ |
| $z$ | 5,000 | $55 \%$ |
| Weighted <br> Average | $2,817.5$ seconds |  |

### 11.3 Mathematical Preliminaries (8 of 17)

- When comparing the relative performance of two or more systems, the geometric mean is the preferred measure of central tendency.
- It is the $n$th root of the product of $n$ measurements.

$$
G=\left(x_{1} \times x_{2} \times x_{3} \times \cdots \times x_{n}\right)^{\frac{1}{n}}
$$

- Unlike the arithmetic means, the geometric mean does not give us a real expectation of system performance. It serves only as a tool for comparison.


## Normalized Radio Using System B as a Reference

- The geometric mean of System A, when using System $B$ as a reference is as below.

$$
\left(\frac{50}{100} \times \frac{200}{400} \times \frac{250}{500} \times \frac{400}{800} \times \frac{5000}{4100}\right)^{\frac{1}{5}}=0.59765
$$

| Program | System A <br> Execution <br> Time | System B <br> Execution <br> Time | System C <br> Execution <br> Time |
| :---: | :---: | :---: | :---: |
| $v$ | 50 | 100 | 500 |
| $w$ | 200 | 400 | 600 |
| $x$ | 250 | 500 | 500 |
| $y$ | 400 | 800 | 800 |
| $z$ | 5,000 | 4,100 | 3,500 |
| Average | 1,180 | 1,180 | 1,180 |

### 11.3 Mathematical Preliminaries (9 of 17)

- The geometric mean is often using normalized ratios between a system under test and a reference machine.
- We have performed the calculation in the table below.

| Program | System A <br> Execution <br> Time | Execution <br> Time <br> Normalized <br> to B | System B <br> Execution <br> Time | Execution <br> Time <br> Normalized <br> to B | System C <br> Execution <br> Time | Execution <br> Time <br> Normalized <br> to B |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $v$ | 50 | 0.5 | 100 | 1 | 500 | 5 |
| $w$ | 200 | 0.5 | 400 | 1 | 600 | 1.5 |
| $x$ | 250 | 0.5 | 500 | 1 | 500 | 1 |
| $y$ | 400 | 0.5 | 800 | 1 | 800 | 1 |
| $z$ | 5,000 | 1.22 | 4,100 | 1 | 3,500 | 0.85366 |
| Geometric <br> Mean |  | 0.59765 |  | 1 |  | 1.44967 |

$$
-\frac{\text { Geometric Mean } A}{\text { Geometric Mean } B}=\frac{0.59765}{1}=0.59765
$$

### 11.3 Mathematical Preliminaries (10 of 17)

- When another system is used for a reference machine, we get a different set of numbers.

| Program | System A <br> Execution <br> Time | Execution <br> Time <br> Normalized <br> to C | System B <br> Execution <br> Time | Execution <br> Time <br> Normalized <br> to C | System C <br> Execution <br> Time | Execution <br> Time <br> Normalized <br> to C |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $v$ | 50 | 0.1 | 100 | 0.2 | 500 | 1 |
| $w$ | 200 | 0.3333 | 400 | 0.66667 | 600 | 1 |
| $x$ | 250 | 0.5 | 500 | 1 | 500 | 1 |
| $y$ | 400 | 0.5 | 800 | 1 | 800 | 1 |
| $z$ | 5,000 | 1.42857 | 4,100 | 1.17143 | 3,500 | 1 |
| Geometric <br> Mean |  | 0.41223 |  | 0.68981 |  | 1 |

- $\frac{\text { Geometric Mean } A}{\text { Geometric Mean } B}=\frac{0.41223}{0.68981}=0.597599$


### 11.3 Mathematical Preliminaries

## (11 of 17)

- The real usefulness of the normalized geometric mean is that no matter which system is used as a reference, the ratio of the geometric means is consistent.
- This is to say that the ratio of the geometric means for System A to System B, System B to System C, and System A to System C is the same no matter which machine is the reference machine.
- Using the results that we got when using System B and System $C$ as reference machines we find that the ratio of the geometric means for System A to System B are $0.59765 / 1=0.41223 / 0.68981$


### 11.3 Mathematical Preliminaries (12 of 17)

- The inherent problem with using the geometric mean to demonstrate machine performance is that all execution times contribute equally to the result.
- So shortening the execution time of a small program by $10 \%$ has the same effect as shortening the execution time of a large program by $10 \%$.
- Shorter programs are generally easier to optimize, but in the real world, we want to shorten the execution time of longer programs.
- Also, if the geometric mean is not proportionate. A system giving a geometric mean $50 \%$ smaller than another is not necessarily twice as fast!


### 11.3 Mathematical Preliminaries

## (13 of 17)

- The harmonic mean provides us with a way to compare execution times that are expressed as a rate, such as operations per second.
- The harmonic mean allows us to form a mathematical expectation of throughput, and to compare the relative throughput of systems and system components.
- To find the harmonic mean, we add the reciprocals of the rates and divide them into the number of rates:

$$
H=n \div\left(1 / x_{1}+1 / x_{2}+1 / x_{3}+\ldots+1 / x_{n}\right)
$$

## Harmonic Mean Example

- Driving 30 miles
- First 10 miles, driving at 30 miles per hour
- Second 10 miles, driving 40 miles per hour
- Last 10 miles, droving 60 miles per hour
- What is average speed
- 30 miles /(time_first10miles + time_second10miles + time_last_10miles)
$=30 /(10 / 30+10 / 40+10 / 60)$
$=3 /(1 / 30+1 / 40+1 / 60)=40$ miles per hour
- Arithmetic average gives wrong result
$-(30+40+60) / 3=43$ miles per hour


### 11.3 Mathematical Preliminaries (14 of 17)

- The harmonic mean holds two advantages over the geometric mean.
- First, it is a suitable predictor of machine behavior.
- So it is useful for more than simply comparing performance.
- Second, the slowest rates have the greatest influence on the result, so improving the slowest performanceusually what we want to do-results in better performance.
- The main disadvantage is that the harmonic mean is sensitive to the choice of a reference machine.


### 11.3 Mathematical Preliminaries (15 of 17)

- This chart summarizes when the use of each of the performance means is appropriate.

| Mean | Uniformly <br> Distributed <br> Data | Data <br> Skewed <br> Data | Dndicator of System <br> Expressed <br> as a Ratio | Data <br> Performance Under <br> a Known Workload | Expressed <br> as a Rate |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Arithmetic | X |  |  | X |  |
| Weighted <br> Arithmetic |  | X |  | X |  |
| Geometric |  | X | X | X | X |
| Harmonic |  |  |  |  |  |

