

MarieSimR: The MARIE Computer Simulator Revision

Xuejun Liang

Department of Computer Science

California State University - Stanislaus

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Introduction

- MarieSim is the MARIE computer simulator
 - MARIE is an accumulator-based computer model used in the popular textbook “The essentials of computer organization and architecture”.
 - Used for assembly language programming exercises.
- But, MarieSim is too simple and thus unable to support some important concepts in computer architecture:
 - No immediate addressing mode
 - No stack and thus its subroutine has no local variables and can not be recursive.
 - Can not define a variable to hold the address of another variable symbolically.
- In order to solve these problems, a revision to MarieSim, called MarieSimR, is developed.

Outline

- Design and Implementation of MarieSimR
- Assembly Language Program Examples
 - to illustrate how to use MarieSimR
 - to compare MarieSim and MarieSimR

The Goal of This Revision

- The goal of this work is to revise MARIE to support the stack and recursive subroutines without changing its architectural characteristics and its microarchitecture for control unit
- The MARIE architecture has the following characteristics:
 - Binary, two's complement data representation.
 - Stored program, fixed word length data and instructions.
 - 4K words of word-addressable main memory.
 - 16-bit data words.
 - 16-bit instructions, 4 for the opcode and 12 for the address.
 - A 16-bit arithmetic logic unit (ALU).
 - Seven registers for control and data movement.

MarieSimR: What's New?

- The stack pointer stored in a reserved memory location and the stack-relative addressing mode are added.
- The subroutine call and return instructions are revised to use the stack for the subroutine return address.
 - A stack frame can be created for a subroutine to hold the return address, input arguments, output results, local variables and so on. So, recursive subroutines are supported.
- A new instruction for increasing or decreasing the value of the stack pointer is added to facilitate the push and pop operations.
- A new instruction for loading an immediate constant into the accumulator is added to replace the clear instruction.
- A new assembler directive is added to support for defining a label to hold the address of another label symbolically

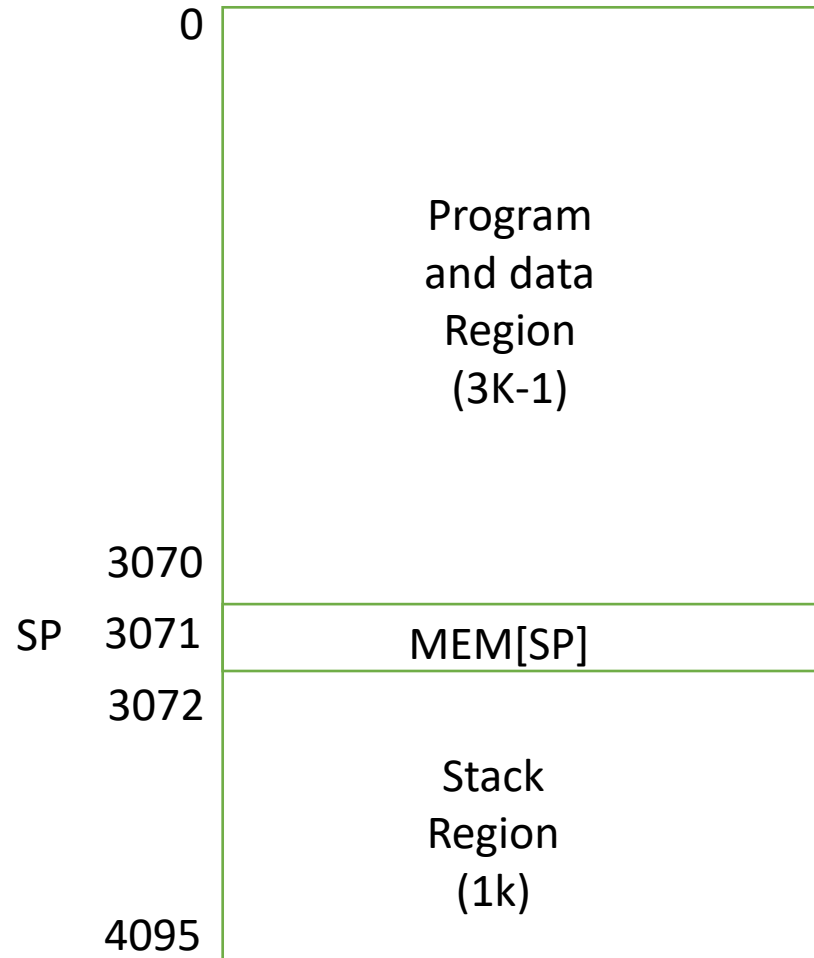
MARIE Instruction Set and Revision

Opcode	Instruction	Meaning
0000	JnS X	$\text{Mem}[X] \leftarrow \text{PC} \ \& \ \text{PC} \leftarrow X+1$
	Call X	Push PC & $\text{PC} \leftarrow X$
0001	Load X	$\text{AC} \leftarrow \text{Mem}[X]$
0010	Store X	$\text{Mem}[X] \leftarrow \text{AC}$
0011	Add X	$\text{AC} \leftarrow \text{AC} + \text{Mem}[X]$
0100	Subt X	$\text{AC} \leftarrow \text{AC} - \text{Mem}[X]$
0101	Input	$\text{AC} \leftarrow$ value from keyboard
0110	Output	Display value in AC on screen
0111	Halt	Terminate program
1000	Skipcond 000	Skip next instruction if $\text{AC} < 0$
	Skipcond 400	Skip next instruction if $\text{AC} = 0$
	Skipcond 800	Skip next instruction if $\text{AC} > 0$
1001	Jump X	$\text{PC} \leftarrow X$
1010	Clear	$\text{AC} \leftarrow 0$
	Limm Imm	$\text{AC} \leftarrow \text{Imm}$
1011	AddI X	$\text{AC} \leftarrow \text{AC} + \text{Mem}[\text{Mem}[X]]$
1100	Jumpl X	$\text{PC} \leftarrow \text{Mem}[X]$
	JR	POP PC
1101	LoadI X	$\text{AC} \leftarrow \text{Mem}[\text{Mem}[X]]$
1110	StoreI X	$\text{Mem}[\text{Mem}[X]] \leftarrow \text{AC}$
1111	IncSP Imm	$\text{Mem}[\text{SP}] += \text{Imm}$

- X is either a hexadecimal literal or a label (symbol) and is used as a memory address
- $\text{Mem}[X]$ represents the content at the memory location X .
- Imm is a 12-bit decimal constant integer.
- When using **stack-relative addressing**, X has the format: $\$ \pm \text{offset}$
 - $\$$ represents the value of the stack pointer
 - **Offset** is a 10-bit decimal constant integer.

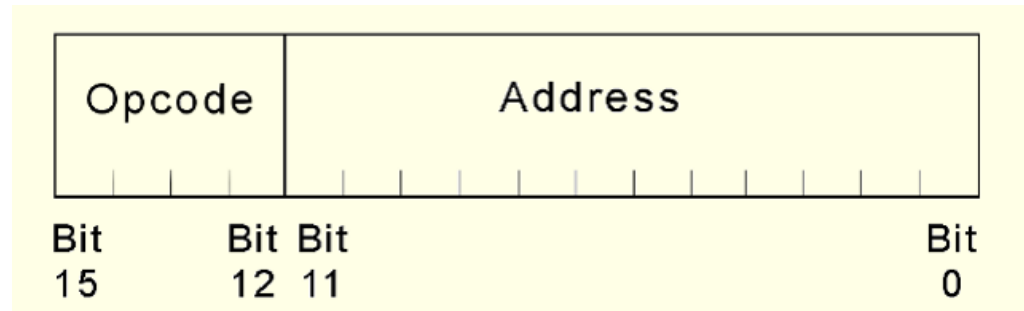
MARIE Memory Map with Stack

- The MARIE architecture has 4K words of word-addressable main memory.
- The stack grows towards the high memory address end and starts from 3072.
- The stack occupies 1K out of 4K MARIE word-addressable memory space.
- The stack pointer is located at memory location 3071

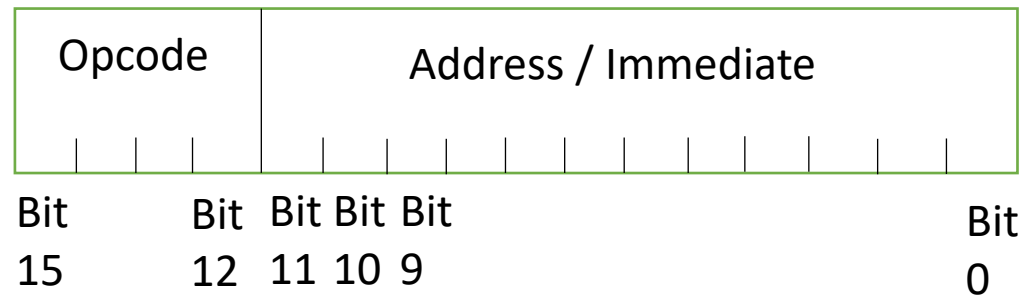


Instruction Encoding Format

MARIE



MARIE
Revision



Instructions **Limm Imm** and **IncSP Imm** store the 12-bit integer in 2's complement inside the instruction from Bit 0 to Bit 11.

If Bit 11 is 1 and Bit 10 is 1, then the stack-relative addressing is used and the offset address is stored inside the instruction from Bit 0 to Bit 9.

Push and Pop Operations

Push X	Meaning
IncSP 1	$\text{Mem}[\text{SP}] \leftarrow \text{Mem}[\text{SP}] + 1$
Load X	$\text{AC} \leftarrow \text{Mem}[X]$
StoreI SP	$\text{Mem}[\text{Mem}[\text{SP}]] \leftarrow \text{AC}$

Pop X	Meaning
loadI SP	$\text{AC} \leftarrow \text{Mem}[\text{Mem}[\text{SP}]]$
Store X	$\text{Mem}[X] \leftarrow \text{AC}$
IncSP -1	$\text{Mem}[\text{SP}] \leftarrow \text{Mem}[\text{SP}] - 1$

The instruction **IncSP Imm** is not necessary as
It can be replaced by three other instructions



IncSP Imm

Limm Imm
Add SP
Store SP

MARIE Directives and Extension

- Directives are instructions to assemblers.
- MarieSim has five directives
 - **ORG** defines the starting address of the program.
 - **DEC**, **OCT**, and **HEX** define named constant in decimal, octa-decimal, and hexadecimal, respectively.
 - **END** indicates the end of the program.
- MarieSimR adds one more directive **LAB**
 - It defines a named hexadecimal constant specified either by a hexadecimal literal or a label symbolically.

Two Programming Examples

- Example 1: Compute sum of numbers in an array.
 - Loop through an array
 - Compare MarieSim and MarieSimR
- Example 2: Compute Fibonacci number Fib(N)
 - Using Loop
 - Using Subroutine with global variables
 - Using Subroutine with local variables
 - Using Recursive subroutine

} Need to use
Stack frame

Example 1: Using loop to add five numbers in array, save result to Sum

```

/ Code for MarieSim and MarieSimR
ORG      100      /Program starts from 0x100
Load    Addr     /Load address of first number to be added
Store   Next     /Store this address is our Next pointer
Load    Num      /Load the number of items to be added
Subt    One      /Decrement
Store   Ctr      /Store this value in Ctr to control looping
Loop,   Load    Sum /Load the Sum into AC
AddI    Next     /Add the value pointed to by location Next
Store   Sum      /Store this sum
Load    Next     /Load Next
Add     One      /Increment by one to point to next address
Store   Next     /Store in our pointer Next
Load    Ctr      /Load the loop control variable
Subt    One      /Subtract one from the loop control variable
Store   Ctr      /Store this new value in loop control variable
Skipcond 000    /If control variable < 0, skip next instruction
Jump    Loop     /Otherwise, go to Loop
Halt

Addr, Hex 117    /Numbers to be summed start at location 117
Next, Hex 0      /A pointer to the next number to add
Num, Dec 5      /The number of values to add
Sum, Dec 0      /The sum
Ctr, Hex 0      /The loop control variable
One, Dec 1      /Used to increment and decrement by 1
Dat, Dec -10    /The values to be added together
Dec 15
Dec -20
Dec 25
Dec 30

```

```

/ Code for MarieSimR
ORG      100
Load    Addr
Store   Next
Limm    -1
Add     Num
Store   Ctr
Loop,   Load    Sum
AddI    Next
Store   Sum
Limm    1
Add     Next
Store   Next
Limm    -1
Add     Ctr
Store   Ctr
Skipcond 000
Jump    Loop
Halt

Addr, Lab  Dat
Next, Hex 0
Num, Dec 5
Sum, Dec 0
Ctr, Hex 0
Dat, Dec -10
Dec 15
Dec -20
Dec 25
Dec 30

```

Example 2: Compute Fibonacci Number fib(N), Where N is an Input

Mathematics formula

$$\text{fib}(N) = \begin{cases} N & \text{if } N < 2 \\ \text{fib}(N - 1) + \text{fib}(N - 2) & \text{if } N \geq 2 \end{cases}$$

C++ code using a loop

```
C = fib(N)
B = fib(N-1)
A = fib(N-2)
```

```
int I, A, B, C, N
cin >> N;
if (N < 2)
    C = N;
else {
    A = 0; B = 1;
    for (I = 2; I <= N; I++) {
        C = B + A; A = B; B = C;
    }
}
cout << C;
```

Method 1 (Using Loop)

```
// Code for MarieSim and MarieSimR
ORG      100
Input    // Read input
Store   N    // N=input
Store   C    // C=N
Subt    I    // AC = N-1
SKIPCOND 800 // if N-1 > 0
JUMP    L3
L2, Load  B    // AC=B
ADD     A    // AC=B+A
Store   C    // C=B+A
Load    B    // AC=B
Store   A    // A=B
Load    C    // AC=C
Store   B    // B=C
Load    I    // AC=1
ADD     One  // AC=I+1
Store   I    // I=I+1
SUBT    N    // AC=I-N
SKIPCOND 400 // if I=N, done
Jump    L2    // otherwise, continue
L3, Load  C    // print result
Output
Halt     // stop
// Variable Declarations
I, DEC   1    // index
N, DEC   0    // N
C, DEC   0    // f(N)
B, DEC   1    // f(N-1)
A, DEC   0    // f(N-2)
One, DEC 1    // Used for increment by 1
```

```
// Code for MarieSimR
ORG      100
Input    // Read input
Store   N    // N=input
Store   C    // C=N
Subt    I    // AC = N-1
SKIPCOND 800 // if N-1 > 0
JUMP    L3
L2, Load  B    // AC=B
ADD     A    // AC=B+A
Store   C    // C=B+A
Load    B    // AC=B
Store   A    // A=B
Load    C    // AC=C
Store   B    // B=C
LIMM    1    // AC=1
ADD     I    // AC=I+1
Store   I    // I=I+1
SUBT    N    // AC=I-N
SKIPCOND 400 // if I=N, done
Jump    L2    // otherwise, continue
L3, Load  C    // print result
Output
Halt     // stop
// Variable Declarations
I, DEC   1    // index
N, DEC   0    // N
C, DEC   0    // f(N)
B, DEC   1    // f(N-1)
A, DEC   0    // f(N-2)
```

Example 2: C++ Code with Using Function

```
int N, C;
int fib(int N) {
    int I, A, B, C;
    if (N < 2)
        C = N;
    else {
        A = 0; B = 1;
        for (I = 2; I <= N; I++) {
            C = B + A; A = B; B = C;
        }
    }
    return C;
}
cin >> N;
C = fib(N);
cout << C;
```

```
int N, C;
int fib(int N) {
    if (N < 2)
        return N;
    else
        return fib(N-1)+fib(N-2);
}
cin >> N;
C = fib(N);
cout << C;
```

Recursive function

Non-recursive function

Method 2 (Using Function and Global Variables)

```
// Code for MarieSim
  ORG 100
  Input          //Read input
  Store N        //N=input
  JNS Fib        //Call subroutine Fib
  Load C         //Print result
  Output
  Halt          //Terminate
// Subroutine Fib
Fib, HEX 0       //for storing return addr
  Load N
  Store C       //C=N
  Subt I        //AC = N-1
  SKIPCOND 800 //if N-1 > 0
  JUMP L3
L2, Load B     //AC=B
  ADD A        //AC=B+A
  Store C      //C=B+A
  Load B       //AC=B
  Store A      //A=B
  Load C       //AC=C
  Store B      //B=C
  Load I       //AC=1
  ADD One      //AC=I+1
  Store I      //I=I+1
  SUBT N       //AC=I-N
  SKIPCOND 400 //if I=N, done
  Jump L2     //otherwise, continue
L3, JumpI Fib
// Global Variable Declarations
I, DEC 1       //index
N, DEC 0       //N -- input to Fib
C, DEC 0       //f(N)-- output from Fib
B, DEC 1       //f(N-1)
A, DEC 0       //f(N-2)
One, DEC 1     //used for increment by 1
```

```
// Code for MarieSimR
  ORG 100
  Input          //Read input
  Store N        //N=input
  Call Fib       //Call subroutine Fib
  Load C         //Print result
  Output
  Halt          //Terminate
// Subroutine Fib
Fib, Load N
  Store C       //C=N
  Subt I        //AC = N-1
  SKIPCOND 800 //if N-1 > 0
  JUMP L3
L2, Load B     //AC=B
  ADD A        //AC=B+A
  Store C      //C=B+A
  Load B       //AC=B
  Store A      //A=B
  Load C       //AC=C
  Store B      //B=C
  Limm 1       //AC=1
  ADD I        //AC=I+1
  Store I      //I=I+1
  SUBT N       //AC=I-N
  SKIPCOND 400 //if I=N, done
  Jump L2     //otherwise, continue
L3, JR
// Global Variable Declarations
I, DEC 1       //index
N, DEC 0       //N -- input to Fib
C, DEC 0       //f(N)-- output from Fib
B, DEC 1       //f(N-1)
A, DEC 0       //f(N-2)
```


Method 3 (Using Function and Local Variables)

Stack frame of Fib(N)

Location in Stack	Memory Address	Used for
\$-1	Mem[SP]-1	Input N / Output Fib(N)
\$	Mem[SP]	Return Address
+\$1	Mem[SP]+1	Local Variable I
+\$2	Mem[SP]+2	Local Variable A
+\$3	Mem[SP]+3	Local Variable B
+\$4	Mem[SP]+4	Local Variable C

```
//Main code for MarieSimR
  ORG      100
  Input           //Read input
  Store    N      //N=input
  Push     N      //Push N on stack (need to expand)
  Call     Fib    //Call subroutine Fib
  Pop      C      //Pop result into C (need to expand)
  Load    C      //print result
  Output
  Halt           //terminate program
```

```

//Subroutine Fib
Fib, Limm 1
    Store $+1 //I=1
    Store $+3 //B=1
    Limm 0
    Store $+2 //A=0
    load $-1 //AC=N
    SUBT $+1 //N-1
    SKIPCOND 800 //if N-1 > 0
    JUMP L3
L2, Load $+3 //AC=B
    ADD $+2 //AC=B+A
    Store $+4 //C=B+A
    Load $+3 //AC=B
    Store $+2 //A=B
    Load $+4 //AC=C
    Store $+3 //B=C
    Limm 1
    ADD $+1 //AC=I+1
    Store $+1 //I=I+1
    SUBT $-1 //AC=I-N
    SKIPCOND 400 //if I=N, done
    Jump L2 //otherwise, continue
    Load $+4 //AC=C
    Store $-1 //Save result
L3, JR
//Global Variable Declarations
N, DEC 0 //N -- input to Fib
C, DEC 0 //f(N)-- output from Fib

```

Initialize the
local variables

Compute Fib(N)

Copy the result to
memory location \$-1

The values in stack at several important timestamps

(1) Right before calling Fib	Memory address	Used for
\$	Mem[SP]	Input N

(2) Right after calling Fib	Memory address	Used for
\$-1	Mem[SP]-1	Input N
\$	Mem[SP]	Return address
+\$1	Mem[SP]+1	Local variable I
+\$2	Mem[SP]+2	Local variable A
+\$3	Mem[SP]+3	Local variable B
+\$4	Mem[SP]+4	Local variable C

(3) Right before Fib return	Memory address	Used for
\$-1	Mem[SP]-1	Output Fib(N)
\$	Mem[SP]	Return address
+\$1	Mem[SP]+1	Local variable I
+\$2	Mem[SP]+2	Local variable A
+\$3	Mem[SP]+3	Local variable B
+\$4	Mem[SP]+4	Local variable C

(4) Right after Fib return	Memory address	Used for
\$	Mem[SP]	Output Fib(N)

Method 4 (Using Recursive Function)

Stack Frame of Fib(N)

Right after calling Fib(N) and right before returning from Fib(N)

Location in Stack	Memory Address	Used for
\$-1	Mem[SP]-1	Input N / Output Fib(N)
\$	Mem[SP]	Return Address
+\$1	Mem[SP]+1	Input N-1 / Output Fib(N-1)
+\$2	Mem[SP]+2	Input N-2 / Output Fib(N-2)

```
//Main code for MarieSimEx
ORG      100
Input                    //Read input
Store     N              //N=input
Push      N              //Push N on stack (need to expand)
JnS       Fib            //Call subroutine Fib
Pop       C              //Pop result into C (need to expand)
Load     C               //print result
Output
Halt                    //terminate program
```

```

//Subroutine Fib
Fib, Limm -1
    Add      $-1      // AC = N-1
    Skipcond 800     // if N > 1
    jump     L1       // Done
    Store    $+1     // Store N-1 to $+1
    IncSP    1       // Increase SP ($) by 1
    Call     Fib     // Call F(N-1)
    Limm     -2      // AC = -2
    Add      $-2     // AC = N-2
    Store    $+1     // Store N-2 to $+1
    IncSP    1       // Increase SP ($) by 1
    Call     Fib     // Call F(N-2)
    IncSP    -2     // Decrease SP ($) by 2 (restore SP)
    Load    $+1     // AC = Fib(N-1)
    Add      $+2     // AC = F(N-1) + F(N-2) = F(N)
    Store    $-1     // store Fib(N) to $-1
L1, JR
//Global Variable Declarations
N, DEC      0       //N -- input to Fib
C, DEC      0       //f(N)-- output from Fib

```

The values in stack at several important timestamps (1)

```
Push N //Push N on stack (need to expand)
```

(1) Right before calling Fib(N)	Memory address	Used for
\$	Mem[SP]	Input N

```
call Fib //Call subroutine Fib
```

(2) Right after calling Fib(N)	Memory address	Used for
\$-1	Mem[SP]-1	Input N
\$+0	Mem[SP]	Return address of Fib(N)
\$+1	Mem[SP]+1	
\$+2	Mem[SP]+2	

```
Store $+1 // Store N-1 to $+1
IncSP 1 // Increase SP ($) by 1
```

(3) Right before calling Fib(N-1)	Memory address	Used for
\$-2	Mem[SP]-2	Input N
\$-1	Mem[SP]-1	Return address of Fib(N)
\$	Mem[SP]	N-1
\$+1	Mem[SP]+1	

The values in stack at several important timestamps (2)

(4) Right after Fib(N-1) return	Memory address	Used for
\$-2	Mem[SP]-2	Input N
\$-1	Mem[SP]-1	Return address
\$	Mem[SP]	Fib(N-1)
+\$1	Mem[SP]+1	

```

Limm    -2    // AC = -2
Add     $-2   // AC = N-2
Store   $+1   // Store N-2 to $+1
IncSP   1     // Increase SP ($) by 1
    
```

(5) Right before calling Fib(N-2)	Memory address	Used for
\$-3	Mem[SP]-3	Input N
\$-2	Mem[SP]-2	Return address
\$-1	Mem[SP]-1	Fib(N-1)
\$	Mem[SP]	N-2

(6) Right after Fib(N-2) return	Memory address	Used for
\$-3	Mem[SP]-3	Input N
\$-2	Mem[SP]-2	Return address
\$-1	Mem[SP]-1	Fib(N-1)
\$	Mem[SP]	Fib(N-2)

The values in stack at several important timestamps (3)

(6) Right after Fib(N-2) return	Memory address	Used for
\$-3	Mem[SP]-3	Input N
\$-2	Mem[SP]-2	Return address
\$-1	Mem[SP]-1	Fib(N-1)
\$	Mem[SP]	Fib(N-2)

```

IncSP    -2    // Decrease SP ($) by 2 (restore SP)
Load     $+1   // AC = Fib(N-1)
Add      $+2   // AC = F(N-1) + F(N-2) = F(N)
Store    $-1   // store Fib(N) to $-1
    
```

(7) Right before Fib(N) return	Memory address	Used for
\$-1	Mem[SP]-1	Fib(N)
\$	Mem[SP]	Return address
+\$1	Mem[SP]+1	Fib(N-1)
+\$2	Mem[SP]+2	Fib(N-2)

```

| L1, JR |
    
```

(8) Right after Fib(N) return	Memory address	Used for
\$	Mem[SP]	Fib(N)

```

| Pop    C    //Pop result into C (need to expand) |
    
```