## More on Solving Recurrence Relations Math 2300, Spring 2017

## Second Order Linear Recurrence Relations with Constant Coefficients

**Theorem**: Let  $p_n$  be ay particular solution to the recurrence relation  $a_k = Aa_{k-1} + Ba_{k-2} + f(k)$ , ignoring initial conditions. Let  $q_n$  be the general solution to the homogeneous recurrence  $a_k = Aa_{k-1} + Ba_{k-2}$  again ignoring initial conditions. Then  $a_n = p_n + q_n$  is the general solution to the recurrence relation  $a_k = Aa_{k-1} + Ba_{k-2} + f(k)$ , where the initial conditions determine the constants in  $q_n$ .

## **Generating Functions:**

A generating function is, roughly, a polynomial that goes on forever:

$$f(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + \dots + a_nx^n + \dots$$

**Definition:** The generating function of a sequence  $a_0$ ,  $a_1$ ,  $a_2$ ,  $a_3$ , ... is the expression  $f(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + ...$ 

$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + \dots$$
 generates 1, 1, 1, ...

$$\frac{1}{1-ax} = 1 + ax + a^2x^2 + a^3x^3 + \dots \text{ generates 1, a, a}^2, a^3, \dots$$

$$\frac{1}{(1-x)^2} = 1 + 2x + 3x^2 + 4x^3 + \dots$$
 generates 1, 2, 3, 4, ...

(The above theorem and definitions adapted from *Discrete Mathematics with Graph Theory* by Goodaire and Parmenter.)

**Recursion Trees** are often used for divide and conquer algorithms.

## The Master Theorem gives bounds on recurrences:

The Master Theorem applies to recurrences of the following form:

$$T(n) = aT(n/b) + f(n)$$

where  $a \ge 1$  and b > 1 are constants and f(n) is an asymptotically positive function.

There are 3 cases:

1. If 
$$f(n) = O(n^{\log_b a - \varepsilon})$$
 for some constant  $\varepsilon > 0$ , then  $T(n) = \Theta(n^{\log_b a})$ .

2. If 
$$f(n) = \Theta(n^{\log_b a} \log^k n)$$
 with  $k \ge 0$ , then  $T(n) = \Theta(n^{\log_b a} \log^{k+1} n)$ .

3. If 
$$f(n) = \Omega(n^{\log_b a + \varepsilon})$$
 with  $\varepsilon > 0$ , and  $f(n)$  satisfies the regularity condition, then  $T(n) = \Theta(f(n))$ . Regularity condition:  $af(n/b) \le cf(n)$  for some constant  $c < 1$  and all sufficiently large n.

For examples of applying the Master Theorem http://www.csanimated.com/animation.php?t=Master\_theorem