First-Order Logic

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Outline

- Why FOL?
- Syntax and semantics of FOL
- Using FOL
- · Wumpus world in FOL
- · Knowledge engineering in FOL

Pros and cons of propositional

- © Propositional logic is declarative
- © Propositional logic allows partial/disjunctive/negated information (unlike most data structures and databases)
- © Propositional logic is compositional:
 - meaning of $B_{1,1} \wedge P_{1,2}$ is derived from meaning of $B_{1,1}$ and of $P_{1,2}$
- Meaning in propositional logic is context-independent
 (unlike natural language, where meaning depends on context)
- (2) Propositional logic has very limited expressive power

 - (unlike natural language)
 E.g., cannot say "pits cause breezes in adjacent squares"
 except by writing one sentence for each square

First-order logic

- Whereas propositional logic assumes the world contains facts,
- first-order logic (like natural language) assumes the world contains
 - Objects: people, houses, numbers, colors, baseball games, wars, ... (squares, pits, wumpuses)
 - Relations: red, round, prime, brother of, bigger than, part of, comes between, ... (is breezy, is adjacent to, shoots)
 - Functions: father of, best friend, one more than, plus, ...

Propositional Logic vs FOL

 $B_{23} \rightarrow (P_{32} \vee P_{23} \vee P_{34} \vee P_{43}) \dots$

"Internal squares adjacent to pits are breezy":

All X Y (B(X,Y) $^{\land}$ (X > 1) $^{\land}$ (Y > 1) $^{\land}$ (Y < 4) $^{\land}$ (X < 4))

 $(P(X-1,Y) \vee P(X,Y-1) \vee P(X+1,Y) \vee (X,Y+1))$

Syntax of FOL: Basic elements

- Constants KingJohn, 2, Pitt,...
- Predicates Brother, >,...
- Functions Sqrt, LeftLegOf,...
- Variables x, y, a, b,...
- Connectives ¬, ⇒, ∧, ∨, ⇔
- Equality
- Quantifiers ∀, ∃

BNF Grammar on p 293

Sentence → AtomicSentence | (Sentence Connective Sentence) | Quantifier Variable, .. Sentence ~Sentence AtomicSentence → Predicate(Term,...) | Term = Term Term → Function(Term,...) | Constant Variable Connective \rightarrow \rightarrow | $^{\wedge}$ | v | \leftarrow \rightarrow Quantifier → all, exists Constant → john, 1, .. Variable \rightarrow A, B, C, X Predicate → breezy, sunny, red Function → fatherOf, plus Knowledge engineering involves deciding what types of things

Detour: Some British History

- Richard the Lionheart
 - Richard I (8 September 1157 6 April 1199) was King of England from 6 July 1189 until his death in 1199.
 - Rebelled unsuccessfully against father Henry II
 - Spoke very little English and mostly lived in Aquitaine
 - Was a central Christian commander during the Third Crusade
 - In his absence, brother John tries to seize throne, Richard forgives him

Atomic sentences

Should be constants, predicates, and functions for your problem

predicate (term₁,...,term_n) Atomic sentence = $\mid term_1 = term_2$

= function (term₁,...,term_n) Term | constant | variable

- Brother(John, Richard)
- Married(Father(Richard), Mother(John))

Complex sentences

• Complex sentences are made from atomic sentences using connectives

$$\neg S$$
, $S_1 \land S_2$, $S_1 \lor S_2$, $S_1 \Rightarrow S_2$, $S_1 \Leftrightarrow S_2$,

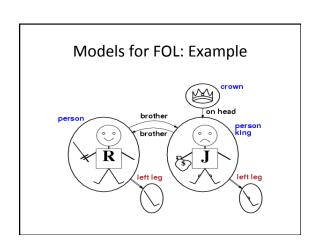
E.g. $Sibling(KingJohn, Richard) \Rightarrow$ Sibling(Richard, KingJohn)

Truth in first-order logic

- Sentences are true with respect to a model and an interpretation
- Model contains objects (domain elements) and relations among them
- Interpretation specifies referents for constant symbols → objects
 predicate symbols → relations
 function symbols → functional relation

Interpretation: assignment of elements from the world to elements of the

An atomic sentence $predicate(term_1,...,term_n)$ is true iff the objects referred to by $term_1,...,term_n$ are in the relation referred to by predicate



Quantifiers

- All X p(X) means that p holds for all elements in the domain
- Exists X p(X) means that p holds for at least one element of the domain

Universal quantification

∀<variables> <sentence>

Everyone at CSU is smart: $\forall x \ At(x,CSU) \Rightarrow Smart(x)$

- $\forall x P$ is true in a model m iff P is true with x being each possible object in the model
- Roughly speaking, equivalent to the conjunction of instantiations of P
 At(KingJohn,CSU) ⇒ Smart(KingJohn)
 At(Richard,CSU) ⇒ Smart(Richard)
 At(CSU,CSU) ⇒ Smart(CSU)

A common mistake to avoid

- Typically, \Rightarrow is the main connective with \forall
- Common mistake: using A as the main connective with ∀:

 $\forall x \text{ At}(x,CSU) \land Smart(x)$ means "Everyone is at CSU and everyone is smart"

Existential quantification

- ∃<variables> <sentence>
- Someone at CSU is smart:
- $\exists x \ At(x,CSU) \land Smart(x)$
- $\exists x \, P \text{ is true} \text{ in a model } m \text{ iff } P \text{ is true} \text{ with } x \text{ being some possible object in the model}$
- Roughly speaking, equivalent to the disjunction of instantiations of P
 - At(KingJohn,CSU) \(\times \) Smart(KingJohn) \(\times \) At(Richard,CSU) \(\times \) Smart(Richard) \(\times \) At(CSU,CSU) \(\times \) Smart(CSU)

Another common mistake to avoid

- Typically, ∧ is the main connective with ∃
- Common mistake: using ⇒ as the main connective with 3:

 $\exists x \, At(x,CSU) \Rightarrow Smart(x)$

is true if there is anyone who is not at CSU!

Transform to: $\exists x \sim (At(x,CSU)) \lor Smart(x)$

Examples

- · Everyone likes chocolate
- · Someone likes chocolate
- · Everyone likes chocolate unless they are allergic to it

Examples

- Everyone likes chocolate
- $\forall X \text{ person}(X) \rightarrow \text{ likes}(X, \text{chocolate})$
- Someone likes chocolate
- ∃X person(X) ^ likes(X, chocolate)
- Everyone likes chocolate unless they are allergic to it ∀X (person(X) ^ ¬allergic (X, chocolate)) → likes(X, chocolate)

Properties of quantifiers

- ∀x ∀y is the same as ∀y ∀x
- $\exists x \exists y \text{ is the same as } \exists y \exists x$
- $\exists x \forall y \text{ is not the same as } \forall y \exists x$
- ∃x ∀y Loves(x,y)
 - "There is a person who loves everyone in the world"
- $\forall y \exists x Loves(x,y)$
 - "Everyone in the world is loved by at least one person"

Nesting of Variables

Put quantifiers in front of likes(P,F) Assume the domain of discourse of P is the set of people Assume the domain of discourse of F is the set of foods

- 1. Everyone likes some kind of food
- 2. There is a kind of food that everyone likes
- 3. Someone likes all kinds of food
- 4. Every food has someone who likes it

Answers (DOD of P is people and F is food)

Everyone likes some kind of food

 \forall P, \exists F likes(P,F)

There is a kind of food that everyone likes

 $\exists F, \forall P \text{ likes}(P,F)$

Someone likes all kinds of food

 $\exists P, \forall F \text{ likes}(P,F)$

Every food has someone who likes it

 \forall F, \exists P likes(P,F)

Answers, without Domain of **Discourse Assumptions**

Everyone likes some kind of food

 \forall P person(P) \rightarrow \exists F food(F) and likes(P,F)

There is a kind of food that everyone likes

 \exists F food(F) and (\forall P person(P) \Rightarrow likes(P,F))

Someone likes all kinds of food

 \exists P person(P) and (\forall F food(F) \Rightarrow likes(P,F))

Every food has someone who likes it

 \forall F food (F) \Rightarrow \exists P person(P) and likes(P,F)

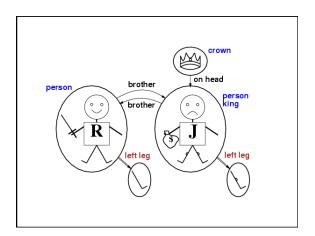
Quantification and Negation

- $\neg \exists x \ p(x) \ equiv \ \forall x \ \neg p(x)$

 - ¬∃x likes(x, parsnips)
 ∀x ¬likes(x, parsnips)
- $\neg \forall x p(x) equiv \exists x \neg p(x)$
 - ¬∀x likes(x, parsnips)
 - ∃x ¬likes(x, parsnips)
- · Quantifier duality: each can be expressed using the other
- $\forall x \text{ Likes}(x, \text{IceCream}) \neg \exists x \neg \text{Likes}(x, \text{IceCream})$
- $\exists x \text{ Likes}(x, \text{Broccoli})$ $\neg \forall x \neg \text{Likes}(x, \text{Broccoli})$

Equality

- term₁ = term₂ is true under a given interpretation if and only if term₁ and term₂ refer to the same object
- E.g., definition of Sibling in terms of Parent:
 ∀x,y Sibling(x,y) ⇔ [¬(x = y) ∧ ∃m,f¬(m = f) ∧ Parent(m,x) ∧ Parent(f,x) ∧ Parent(m,y) ∧ Parent(f,y)]



- Predicate of brotherhood:
 - {<R,J>,<J,R>}
- Predicate of being on: {<C,J>}
- Predicate of being a person:
 - {J,R}
- Predicate of being the king: {J}
- Predicate of being a crown: {C}
- Function for left legs: <{J,JLL},{R,RLL}>

Interpretation

- Specifies which objects, functions, and predicates are referred to by which constant symbols, function symbols, and predicate symbols.
- Under the intended interpretation:
- "richardl" refers to R; "johnII" refers to J; "crown" refers to the crown.
- "onHead","brother","person","king", "crown", "leftLeg", "strong"

Lots of other possible interpretations

- 5 objects, so just for constants "richard" and "john" there are 25 possibilities
- Note that the legs don't have their own names!
- "johnII" and "johnLackland" may be assigned the same object, J
- Also possible: "crown" and "john!!" refer to C (just not the intended interpretation)

Why isn't the "intended interpretation" enough?

- Vague notion. What is intended may be ambiguous (and often is, for non-toy domains)
- Logically possible: square(x) ^ round(x). Your KB has to include knowledge that rules this out.

Determining truth values of FOPC sentences

- Assign meanings to terms:
 - "johnII" ← J; "leftLeg(johnII)"← JLL
- · Assign truth values to atomic sentences
 - "brother(johnII,richardI)"
 - "brother(johnlackland,richardI)"
 - Both True, because <J,R> is in the set assigned "brother"
 - "strong(leftleg(johnlackland))"
 - True, because JLL is in the set assigned "strong"

Examples given the Sample Interpretation

- ∀ X,Y brother(X,Y) FALSE
- ∀ X,Y (person(X) ^ person(Y)) → brother(X,Y) FALSE
- ∀ X,Y (person(X) ^ person(Y) ^ ~(X=Y)) → brother(X,Y) TRUE
- ∃ X crown(X) TRUE
- ∃ X ∃ Y sister(X,Y) FALSE

Representational Schemes

- What are the objects, predicates, and functions?
 Keep in mind that you need to encode knowledge of specific problem instances and general knowledge.
- In practice, consider interpretations just to understand what the choices are. The world and interpretation are defined, or at least constrained, through the logical sentences we write.

Example Choice: Predicates versus Constants

Rep-Scheme 1: Let's consider the world: D = {a,b,c,d,e}. green: {a,b,c}. blue: {d,e}. Some sentences that are satisfied by the intended interpretation:

green(a). green(b). blue(d). $\sim (\forall x \text{ green}(x)). \forall x \text{ green}(x) \text{ v blue}(x).$

But what if we want to say that blue is pretty?

Choice: Predicates versus Constants

- Rep-Scheme 2: The world: D = {a,b,c,d,e,green,blue} colorof: {<a,green>,<b,green>,<c,green>,<d,blue>,<e,blue>} pretty: {blue} notprimary: {green}
- Some sentences that are satisfied by the intended interpretation: colorOf(a,green). colorOf(b,green). colorOf(d,blue).
 ~(∀ x colorOf(x,green)).

▼ X colorOf(X,green) v colorOf(X,blue).
pretty(blue). notprimary(green).

pretty(blue). notprimary(green).
We have reified predicates blue and green: made them into objects

Using FOL

The kinship domain:

- Brothers are siblings
 ∀x,y Brother(x,y) ⇔ Sibling(x,y)
- One's mother is one's female parent
 ∀m,c Mother(c) = m ⇔ (Female(m) ∧ Parent(m,c))
- "Sibling" is symmetric
 ∀x,y Sibling(x,y) ⇔ Sibling(y,x)

Interacting with FOL KBs

Suppose a wumpus-world agent is using an FOL KB and perceives a smell and a breeze (but no glitter) at t=5:

Tell(KB,Percept([Smell,Breeze,None],5)) Ask(KB, 3a BestAction(a,5))

- I.e., does the KB entail some best action at t=5?
- Answer: Yes, {a/Shoot} ← substitution (binding list)
- Given a sentence S and a substitution σ,
- $\emph{S}\sigma$ denotes the result of plugging σ into $\emph{S};$ e.g.,
 - S = Smarter(x,y)
- σ = {x/Hillary,y/Bill} Sσ = Smarter(Hillary,Bill)
- As k(KB,S) returns some/all σ such that KB \vdash S σ

Knowledge base for the wumpus world

- Perception
 - \forall t,s,b Percept([s,b,Glitter],t) \Rightarrow Glitter(t)
- - ∀t Glitter(t) \Rightarrow BestAction(Grab,t)

Deducing hidden properties

• $\forall x,y,a,b \ Adjacent([x,y],[a,b]) \Leftrightarrow$ $[a,b] \in \{[x+1,y], [x-1,y], [x,y+1], [x,y-1]\}$

Properties of squares:

∀s,t At(Agent,s,t) ∧ Breeze(t) ⇒ Breezy(s)

Squares are breezy near a pit:

- Diagnostic rule---infer cause from effect
- $\forall s \text{ Breezy}(s) \Rightarrow \exists r, Adjacent(r,s) \land Pit(r)$
- Causal rule---infer effect from cause
 - $\forall r \; \mathsf{Pit}(r) \Rightarrow [\forall s \; \mathsf{Adjacent}(r,s) \Rightarrow \mathsf{Breezy}(s) \;]$

Knowledge engineering in FOL

- 1. Identify the task
- 2. Assemble the relevant knowledge
- 3. Decide on a vocabulary of predicates, functions, and
- 4. Encode general knowledge about the domain
- 5. Encode a description of the specific problem instance
- 6. Pose queries to the inference procedure and get answers
- 7. Debug the knowledge base

Summary

- First-order logic:
 - objects and relations are semantic primitives
 - syntax: constants, functions, predicates, equality, quantifiers
- Increased expressive power: better to define wumpus world