

Adversarial Search

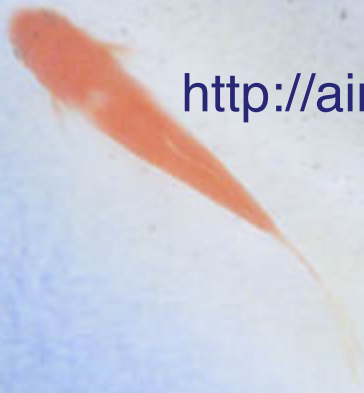
Dr. Melanie Martin

CS 4480

October 1, 2010

Based on slides from

<http://aima.eecs.berkeley.edu/2nd-ed/slides-ppt/>



Outline

- Optimal decisions
- α - β pruning
- Imperfect, real-time decisions

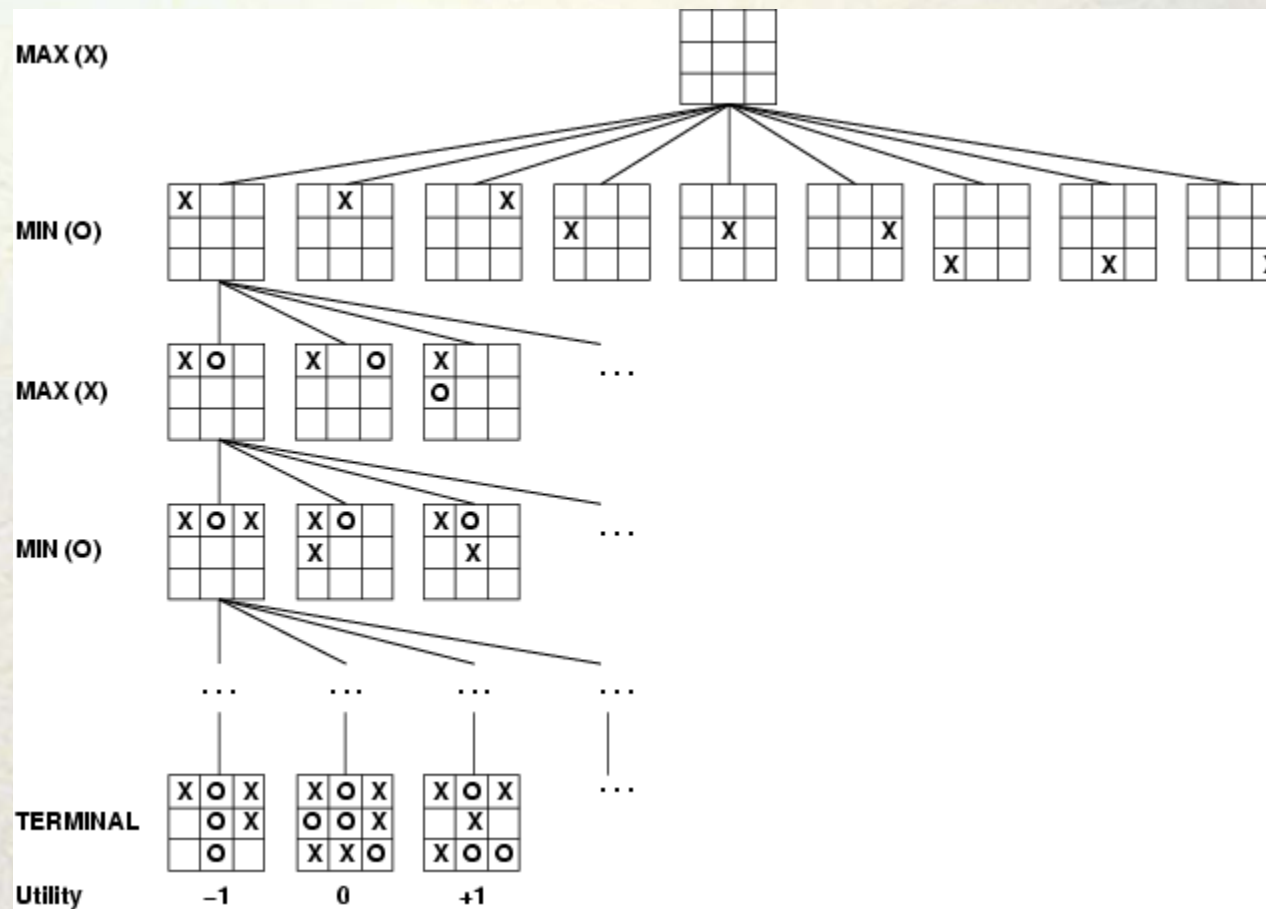
Games vs. search problems

- "Unpredictable" opponent → specifying a move for every possible opponent reply
- Time limits → unlikely to find goal, must approximate

Game as Search Problem

- Initial State
 - Board position and player to move
- Successor Function
 - Returns list of *(state, move)* pairs
 - Legal moves and resulting states
- Terminal (Goal) Test
 - When game is over
- Utility (Objective) Function
 - Assigns numeric outcome to terminal states
 - E.g. +1, -1, 0 for win, lose, draw

Game tree (2-player, deterministic, turns)

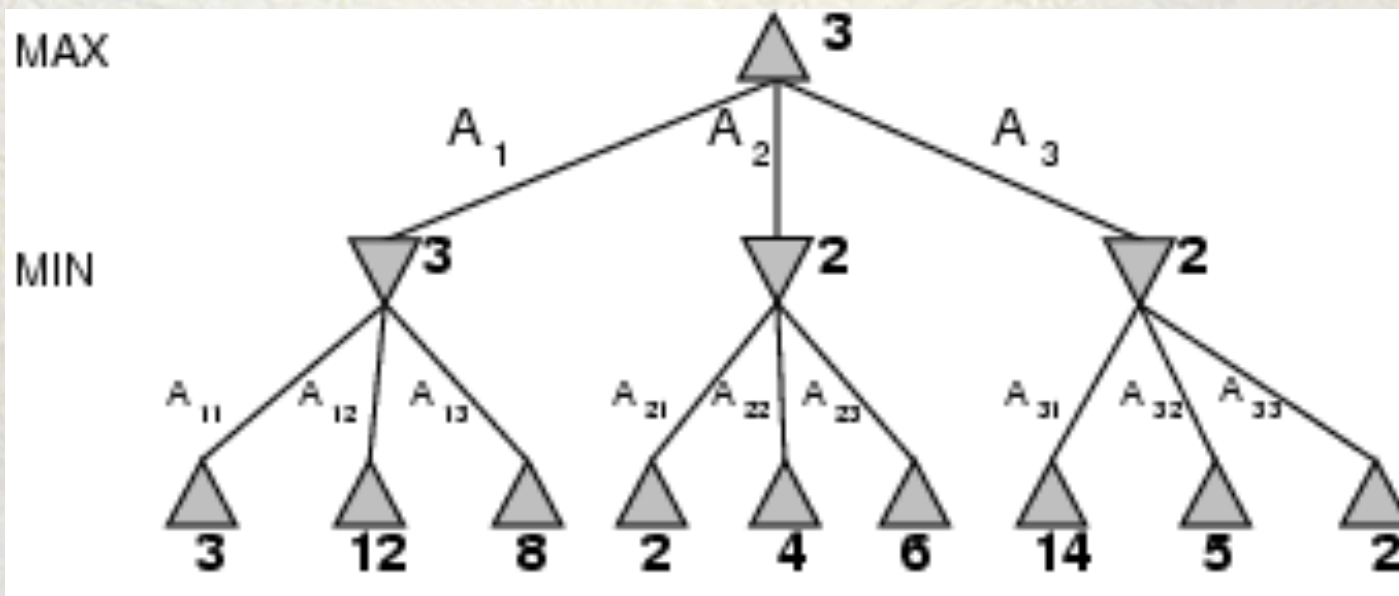


Optimal Strategy

- Leads to outcomes as good as any other strategy when playing an infallible opponent
- Tree where max takes a turn and min takes a turn is ONE MOVE DEEP made up two half-moves - each half move is called a ply

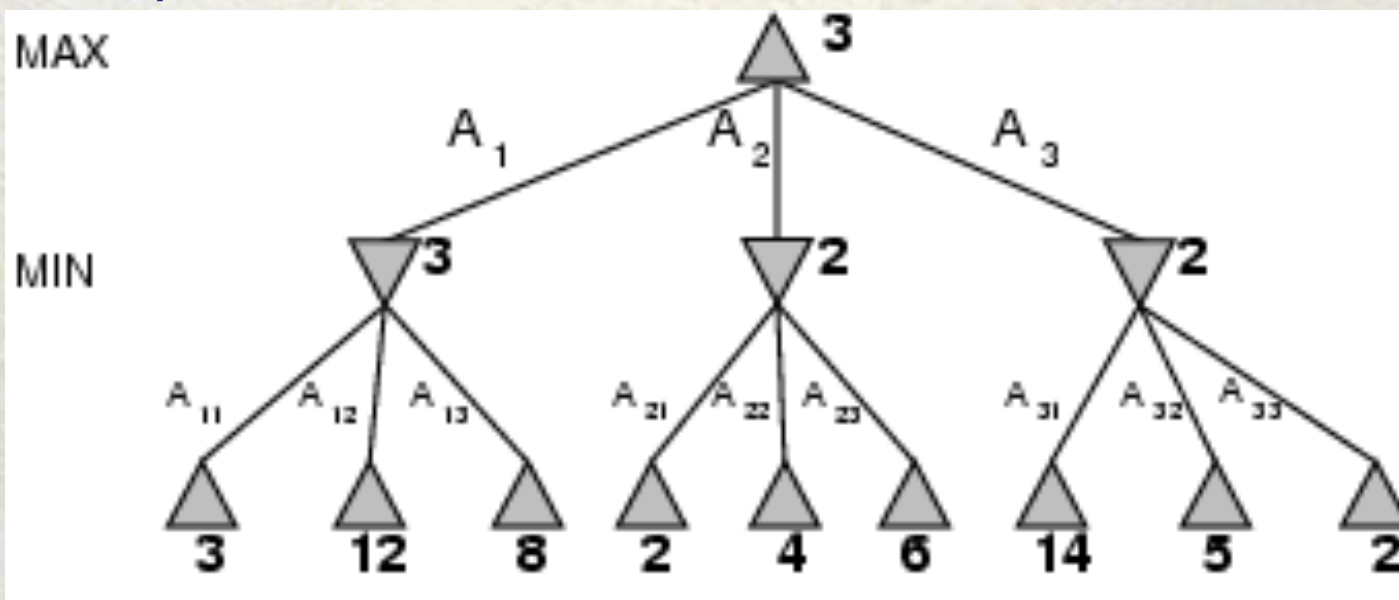
Minimax

- Perfect play for deterministic games
- Idea: choose move to position with highest **minimax value** = best achievable payoff against best play
- E.g., 2-ply game:



Minimax

- Minimax(node) is utility for max of being in corresponding state
- Max prefers a state with maximum value
- Min prefers a state with minimum value



Minimax algorithm

function MINIMAX-DECISION(*state*) *returns an action*

$v \leftarrow \text{MAX-VALUE}(\textit{state})$

return the *action* in SUCCESSORS(*state*) with value *v*

function MAX-VALUE(*state*) *returns a utility value*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow -\infty$

for *a, s* in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s))$

return *v*

function MIN-VALUE(*state*) *returns a utility value*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow \infty$

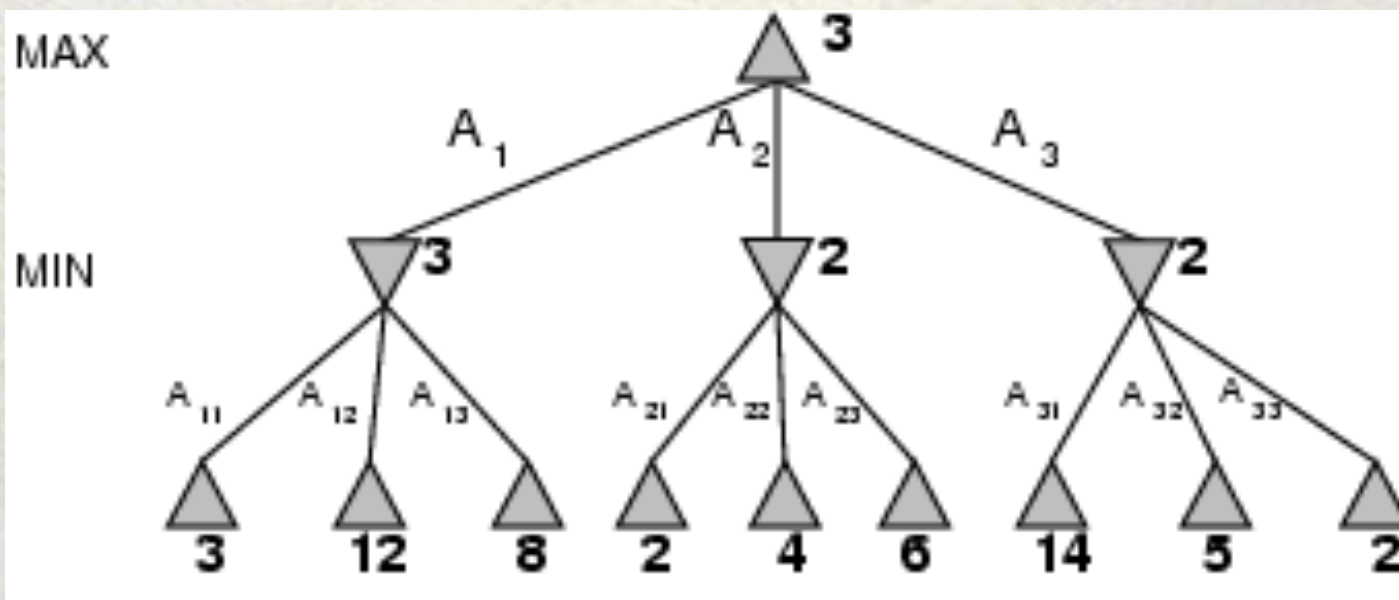
for *a, s* in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s))$

return *v*

Minimax

- Recursion proceeds to leaves and based on utility function assigns minimax values at the level above and so on



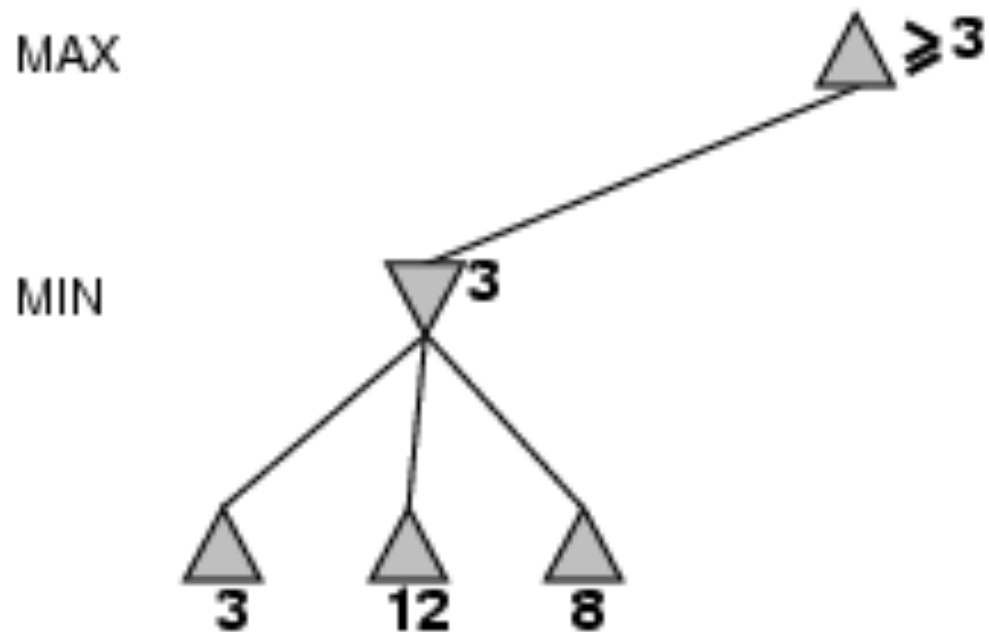
Properties of minimax

- Complete? Yes (if tree is finite)
- Optimal? Yes (against an optimal opponent)
- Time complexity? $O(b^m)$
- Space complexity? $O(bm)$ (depth-first exploration)
- For chess, $b \approx 35$, $m \approx 100$ for "reasonable" games
→ exact solution completely infeasible

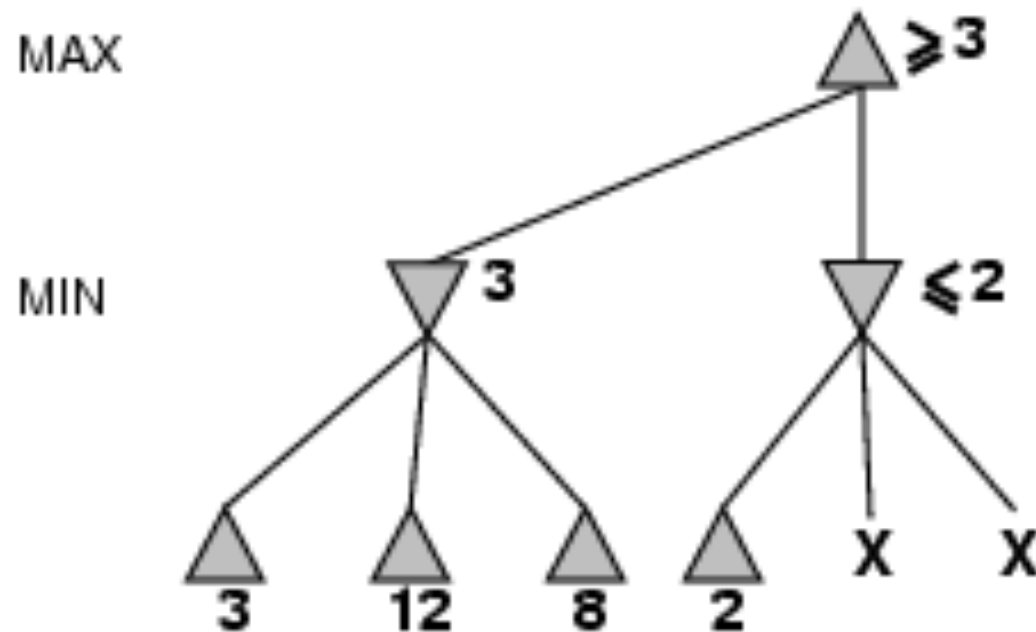
Problem

- Number of game states exponential in number of moves
- Can cut in half with pruning
 - Still exponential
 - Get rid of branches that can't influence final decision

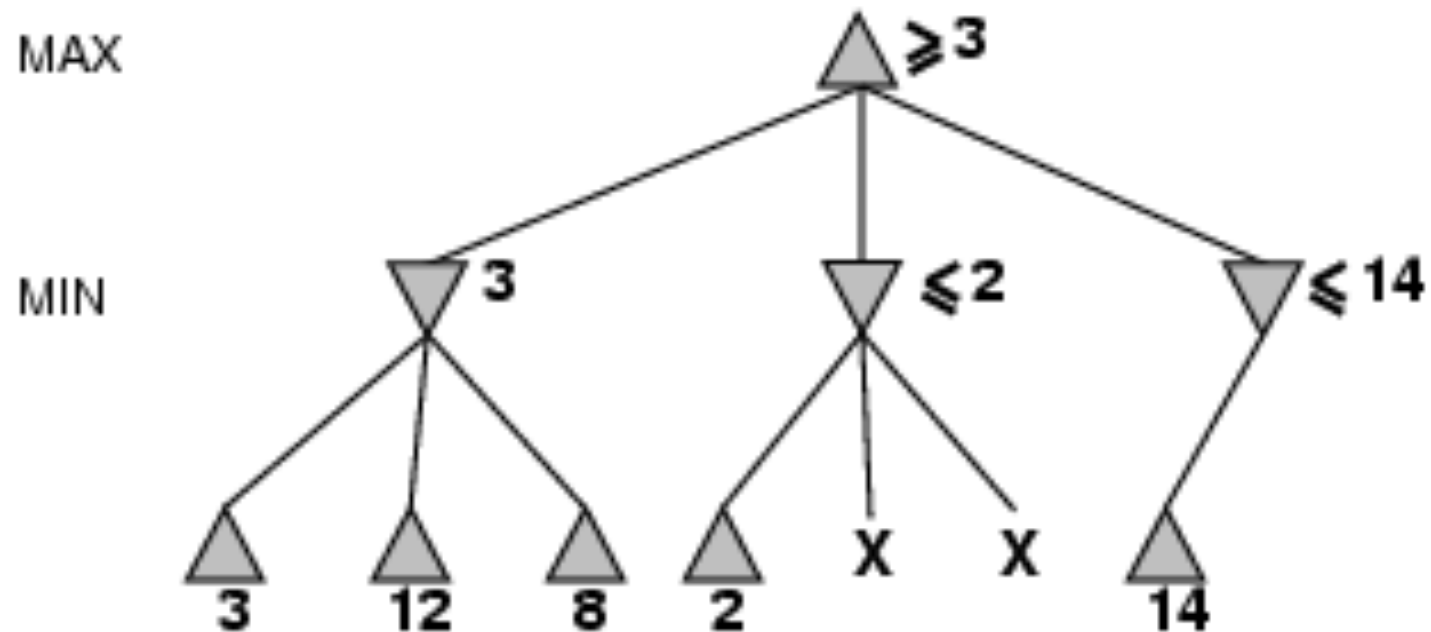
α - β pruning example



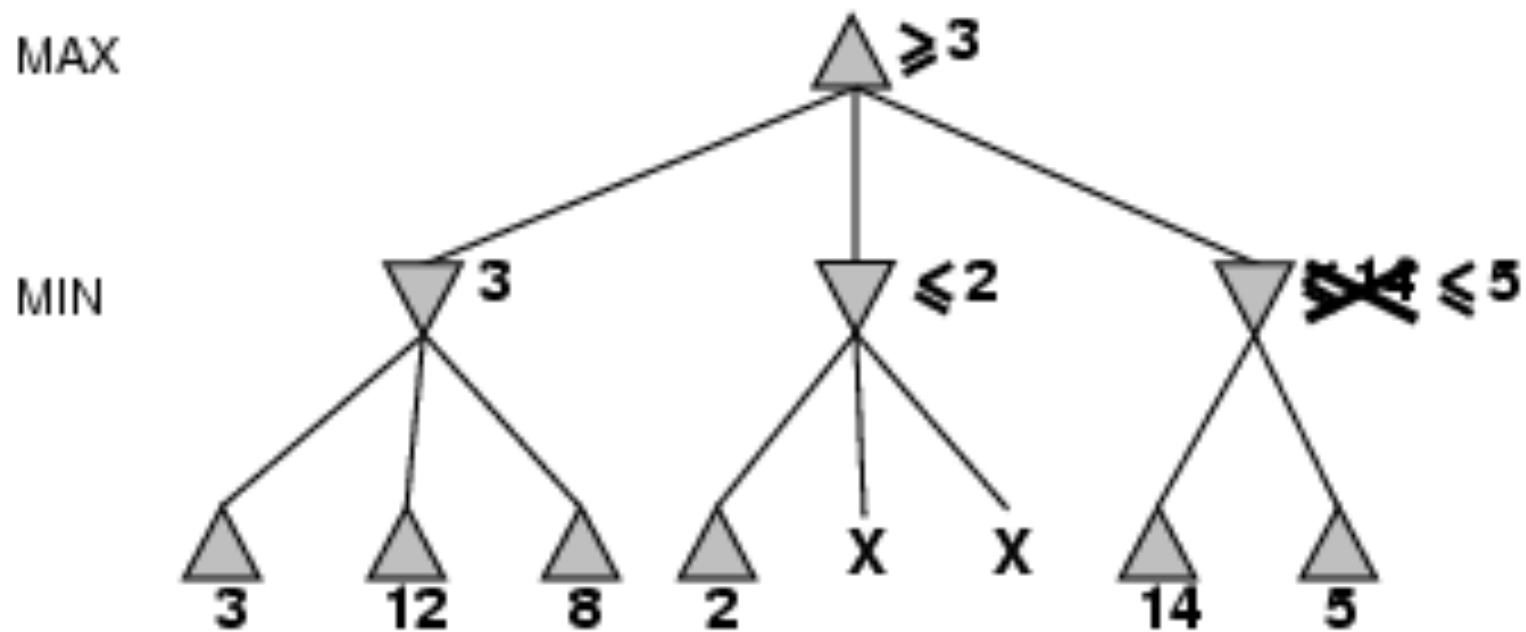
α - β pruning example



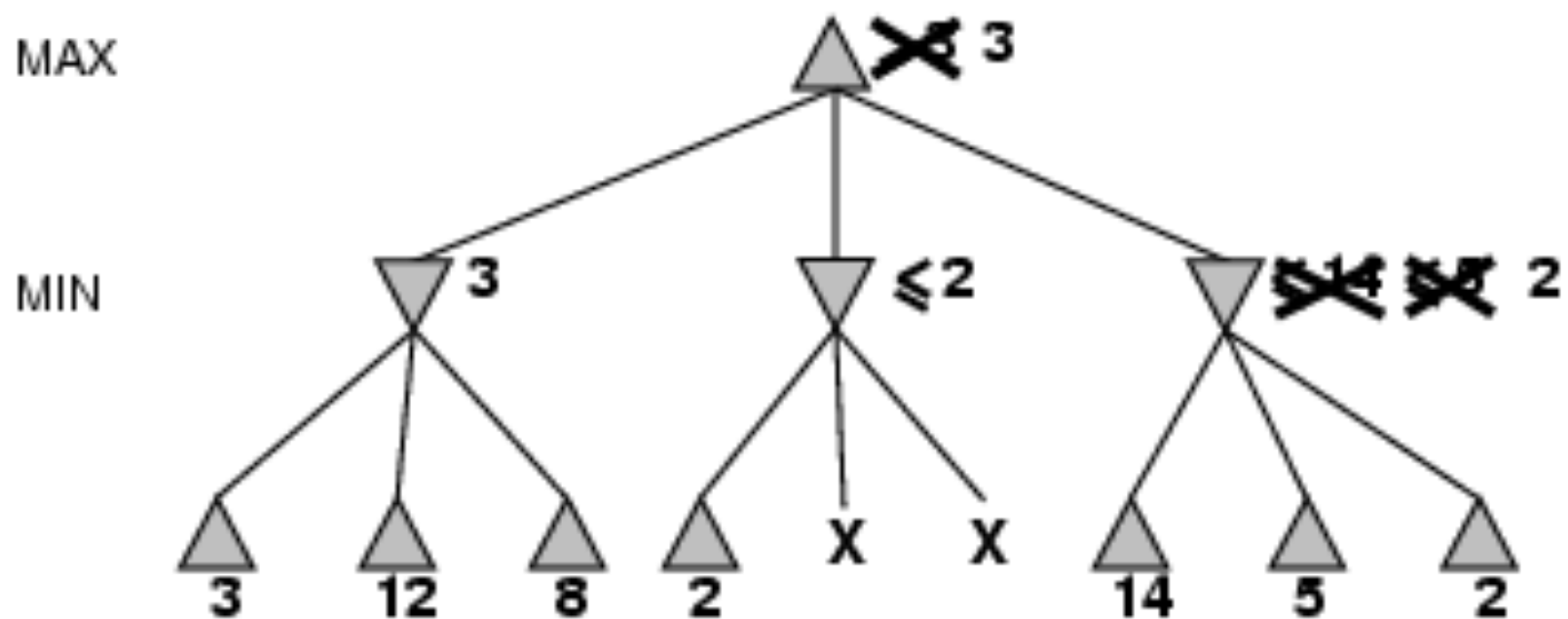
α - β pruning example



α - β pruning example



α - β pruning example

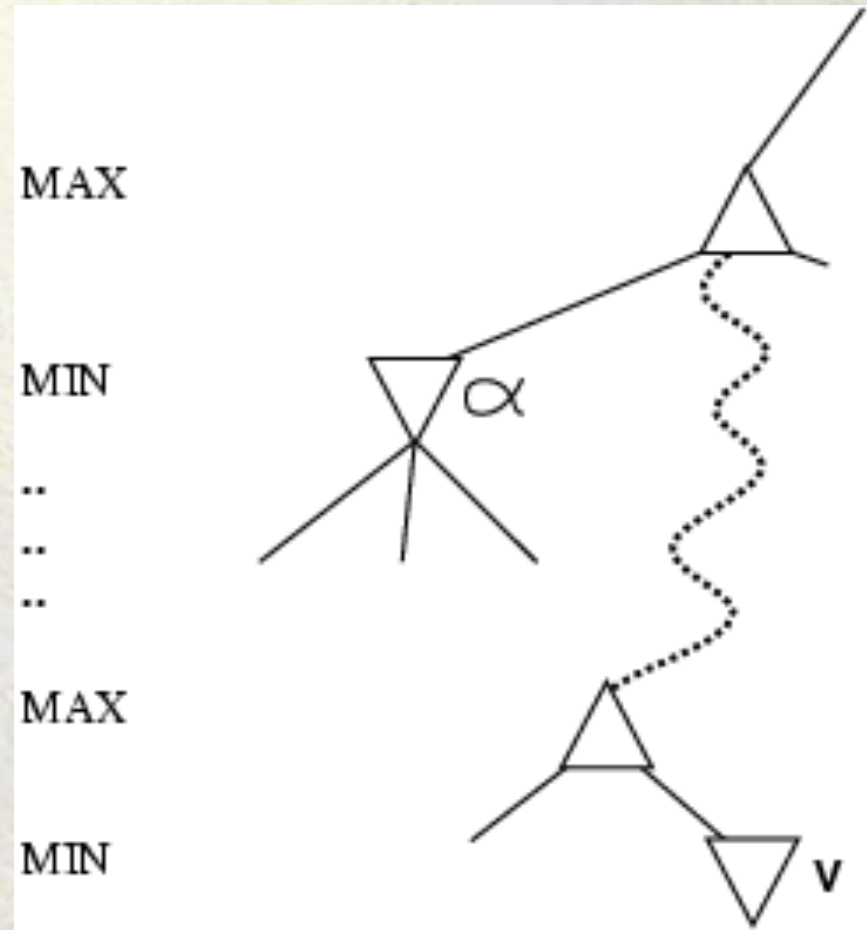


Properties of α - β

- Pruning **does not** affect final result
- Good move ordering improves effectiveness of pruning
- With "perfect ordering," time complexity = $O(b^{m/2})$
→ **doubles** depth of search
- A simple example of the value of reasoning about which computations are relevant (a form of **metareasoning**)

Why is it called α - β ?

- α is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for *max*
- If v is worse than α , *max* will avoid it
→ prune that branch
- Define β similarly for *min*



The α - β algorithm

function ALPHA-BETA-SEARCH(*state*) *returns an action*

inputs: *state*, current state in game

$v \leftarrow \text{MAX-VALUE}(\text{state}, -\infty, +\infty)$

return the *action* in SUCCESSORS(*state*) with value v

function MAX-VALUE(*state*, α , β) *returns a utility value*

inputs: *state*, current state in game

α , the value of the best alternative for MAX along the path to *state*

β , the value of the best alternative for MIN along the path to *state*

if TERMINAL-TEST(*state*) **then return** UTILITY(*state*)

$v \leftarrow -\infty$

for a, s in SUCCESSORS(*state*) **do**

$v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(s, \alpha, \beta))$

if $v \geq \beta$ **then return** v

$\alpha \leftarrow \text{MAX}(\alpha, v)$

return v

The α - β algorithm

function MIN-VALUE($state, \alpha, \beta$) *returns a utility value*

inputs: $state$, current state in game

α , the value of the best alternative for MAX along the path to $state$

β , the value of the best alternative for MIN along the path to $state$

if TERMINAL-TEST($state$) **then return** UTILITY($state$)

$v \leftarrow +\infty$

for a, s in SUCCESSORS($state$) **do**

$v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(s, \alpha, \beta))$

if $v \leq \alpha$ **then return** v

$\beta \leftarrow \text{MIN}(\beta, v)$

return v

Resource limits

Suppose we have 100 secs, explore 10^4 nodes/sec
→ 10^6 nodes per move

Standard approach:

- **cutoff test:**
e.g., depth limit
- **evaluation function**
= estimated desirability of position
Usually a heuristic function

Evaluation functions

- For chess, typically **linear** weighted sum of **features**

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

- e.g., $w_1 = 9$ with
 $f_1(s) = (\text{number of white queens}) - (\text{number of black queens}), \text{ etc.}$

Cutting off search

MinimaxCutoff is identical to *MinimaxValue* except

1. *Terminal?* is replaced by *Cutoff?*
2. *Utility* is replaced by *Eval*

Does it work in practice?

$$b^m = 10^6, b=35 \rightarrow m=4$$

4-ply lookahead is a hopeless chess player!

- 4-ply $\approx$ human novice
- 8-ply $\approx$ typical PC, human master
- 12-ply $\approx$ Deep Blue, Kasparov

Deterministic games in practice

- Checkers: Chinook ended 40-year-reign of human world champion Marion Tinsley in 1994. Used a precomputed endgame database defining perfect play for all positions involving 8 or fewer pieces on the board, a total of 444 billion positions.
»
- Chess: Deep Blue defeated human world champion Garry Kasparov in a six-game match in 1997. Deep Blue searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply.
- Othello: human champions refuse to compete against computers, who are too good.
- Go: human champions refuse to compete against computers, who are too bad. In go, $b > 300$, so most programs use pattern knowledge bases to suggest plausible moves.

Summary

- Games are fun to work on!
- They illustrate several important points about AI
- perfection is unattainable → must approximate
- good idea to think about what to think about