Algol Part 1

CS4100
February 24, 2012

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After FORTRAN

- International language is needed
  - 1964: New language is proposed to break away from platform dependence
  - Preliminary spec: NPL (new programming language), then PL/I (programming language 1)
  - PL/I is too big
    - Dijkstra: If Fortran is an infantile disorder, then PL/I is a fatal disease
    - Trying to be everything to everyone backfires

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Chapter 3: Generality and Hierarchy: ALGOL-60

- An international language is needed
  - A single, universal language would be valuable
  - International (American and European) committee is set up to make recommendations
  - Algol-58 is created in 8 days in Zurich, as a preliminary report
  - Algol: Algorithmic Language

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Implementations

- Because of the hype, many started implementation quickly
  - This resulted in many dialects
  - JOVIAL (Jules' Own Version of the International Algebraic Language)
- Committee meets again in 1960 to incorporate suggestions
  - Algol-60 is born and is very different from the '58 report.
  - Report is 17 pages long: remarkable achievement, mainly due to BNF notation (reports used to stretch to hundreds or thousands of pages)

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Algol Report

- 1959 UNESCO Conference on Information Processing
  - Backus presents a description of Algol '58
    - Uses formal syntax he developed
  - Naur is editor of Algol Bulletin
    - Disagrees with some of Backus' interpretation
    - Need for more precise description
    - Develops a variant of Backus' formal syntax

Backus-Naur Form, aka BNF used for 1960 Algol Report

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Algol's Objectives

- The language should be very close to mathematical notation
- Should be useful in publications to describe algorithms
- Mechanically translatable to machine code

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Structural Organization

- Hierarchically structured language
  - Nesting is introduced (Fortran did not use nesting)
  - Control structures can also be nested
    - One can be made the body of the other

```
if N > 0 then
  for i := 1 step 1 until N do
        sum := sum + Data[i]
```

    - Advantage: decreases the number of GOTOs required
- Reserved words

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Constructs

- Declarative or Imperative
  - (like in FORTRAN)
- Variable declarations
  - Types: integer, real, Boolean

```
integer i, j, k
```

  - Lower bounds of arrays need not be 1

```
real array Data[-50:50]
```

  - Switch, like FORTRAN's computed GOTO
- Subprogram declarations
  - Keyword: **procedure** and
  - Procedures can be *typed* (functions) and *untyped*

```
real procedure dist(x1, y1, x2, y2);
  real x1, y1, x2, y2;
  dist = sqrt((x1 - x2)^2 + (y1 - y2)^2)
```

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Imperative Constructs

- Computational
  - Assignment: "variable := expression"
  - Operators:
    - Arithmetic: +, -, *, etc.
    - Relational: =, <, >, ≥, etc.
    - Logic: ∧, ∨, ¬, etc.
  - Why is assignment ':=' and not '='?
    - Assignment is different from definition and comparison
    - I = I + 1 ; I := I + 1

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Imperative Constructs

- Control-flow
  - All imperative constructs alter flow of control (except assignment)
  - Has *if-then-else*
  - *for*-loop replaces *do*-loop
- No input/output constructs
  - I/O was left to be handled by platform-dependent library calls

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Name Structures

- Algol-60 introduces the compound statement
  - Where 1 statement is allowed, more can be used, using begin-end

```
for i := 1 step 1 until N do
      sum := sum + Data[i]

for i := 1 step 1 until N do
  begin
      sum := sum + Data[i];
      Print Real (sum)
  end
```

  - Also, the body of a procedure is a single statement

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Syntax - Program

- <program> ::= <block> | <compound statement>
- <block> ::= <unlabelled block> | <label>: <block>
- <compound statement> ::= <unlabelled compound> | <label>: <compound statement>
- <unlabelled compound> ::= begin <compound tail>
- <unlabelled block> ::= <block head> ; <compound tail>

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Syntax - Block

- <block> ::= <unlabelled block> |
  <label>: <block>
- <unlabelled block> ::=
  <block head> ; <compound tail>
- <block head> ::= begin <declaration> |
  <block head> ; <declaration>
- <compound tail> ::= <statement> end |
  <statement> ; <compound tail>

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Syntax - Statement

- <compound statement> ::= <unlabelled compound> |
  <label>: <compound statement>
- <unlabelled compound> ::= begin <compound tail>
- <compound tail> ::= <statement> end | <statement> ; <compound tail>
- <statement> ::= <unconditional statement> | <conditional statement> |
  <for statement>
- <unconditional statement> ::= <basic statement> |
  <compound statement> | <block>
- <basic statement> ::= <unlabelled basic statement> |
  <label>: <basic statement>
- <unlabelled basic statement> ::= <assignment statement> |
  <go to statement> | <dummy statement> |
  <procedure statement>

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Name Binding

- Fortran binds a variable to a single memory location statically
- Algol-60 included the results of research done on name structures, which were problematic in Fortran
  - Sharing of data between subprograms
  - Parameter passing modes
  - Return values
  - Dynamic arrays
- Result of research: block structure

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Blocks Define Nested Scopes

- Fortran
  - Had local and global declarations only
  - Had to redeclare using COMMON to share
- Algol-60
  - Introduces blocks

```
begin
  declarations;
  statements
end
```

  - Compound statements do not have '`declarations`'.
  - All declarations are visible to all statements in the block
  - Since statements can be blocks, scopes can be nested

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Why do we need scopes?

- Reduce the context programmers have to keep in mind
- Make understanding and maintenance of program easier
- Scopes reduce visibility of names
  - Declare variable only where needed and used
- Nested blocks inherit names from outside
  - It would be very inconvenient if they did not

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"COMMON" with Blocks

- The error-prone COMMON in Fortran can be implemented in a much better way using blocks

```
begin
  integer array Name, Loc, Type[1:100];
  procedure Lookup (n);
          . . . Lookup procedure . . .
  procedure Var (n, l, t);
          . . . Var procedure . . .
  procedure Array1 (n, l, t, dim1);
          . . . Array1 procedure . . .
end
```

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Too Much Access

- Blocks provide “indiscriminate access”
  - Since functions must be accessible to users,
  - and data structures must be accessible to functions
  - → Data is also accessible to users
- Violates information hiding principle

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Contour Diagrams

- Inner blocks implicitly inherit access to all variable in immediately surrounding block
- Names declared in a block are local to the block
- Names declared in surrounding blocks are nonlocal
- Names declared in outermost block are global

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