

## Set Types

- Pascal provides facilities for sets
set of <ordinal type>
- Ordinal type: enumeration, char, Boolean, subrange
- Not integer or real
var $S, T:$ set of $1 . .10$;
- S, T can hold a set of numbers between 1 and 10 - vs a single number between 1 and 10: var $S, T: 1 . .10 ;$


## Efficiency of Sets

- Set types are restricted to be ordinal to be efficient
var $S, T:$ set of $1 . .10$;
-S, T take only 10 bits to represent: 1 bit for each number
- Bit $=0$ means number is not is set
- Bit = 1 means number is in set
$-S:=[1,2,3,5,7]$;

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ | $\mathbf{8}$ | $\mathbf{9}$ | $\mathbf{1 0}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathrm{S}=$ | 1 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 0 |

## Set Operations

- Initialization/Assignment

T : = [1..6];

- Membership
in
if 4 in $T$ then
- Union, intersection, difference
,+ *,
S * T, $\mathbf{S}+\mathrm{T}$,
- Comparisons
- Subset, equality, non-equality
- <=, >=, =, <>
- Proper subset (<) is not provided


## Efficiency of Sets

- Sets are implemented using bit masks
- Therefore, operations on sets can be implemented using logical operations
- Intersection: logical and
- Union: logical or
- Difference: logical exclusive or
- Logical operations are the fastest a computer can do
- Memory efficiency: 1 bit per element


## Sets

- Considered an example of elegance
- High-level
- Readable
- Efficient
- Secure



## Array Types

- Arrays are more general than Algol's
- Base type of arrays can be non-primitives
- Index types are introduced
- Subscripts can be other than integers
- Char, subrange, enumerated types
var A: array [1..100] of real;
var Occur: array [char] of integer;
var HoursWorked: array [Mon..Fri] of 0..24;
for day := Mon to Fri do
TotalHours := TotalHours + HoursWorked[day];


## Dimensions

- Only single-dimension arrays are allowed!!!
- However:
- Base type of array can be another array!!! var M: array [1..20] of array [1..100] of real;
- Dereferencing: m[3] [5]
- Syntactic sugar.
var M: array [1..20, 1..100] of real; M[3, 5]
(Doesn't affect functionality, sweeter for human use.)


## Static Arrays Only

- Algol's dynamic arrays are not supported
- Type checking is done at compile time
- Array bounds are part of array type
- Hence, only static arrays are supported


## Record Types

- Pascal provides the ability to group heterogeneous data
- Versus homogeneous, using arrays
- Can contain any other type, even other records type person $=$ record


## name: string;

age: 16..100;
salary: 10000..100000;
sex: (male, female);
hireDate: date;
end;
string $=$ array [1..30] of char;

## Dereferencing Records

- Dereferencing is done using the '.
var today: date;
newhire.age := 25
newhire.hireDate := today;
newhire.hireDate.month $:=$ Mar; if newhire.name[1] = ' $A$ ' then ..
employee[en]. hireDate.year $:=2004$;
- Opening one record for multiple access
with newhire do
begin
age := 25;
hireDate := today;
hireDate.month $:=$ Mar;
end;


## Variant Records

- Pascal supports saving storage using variant records; allows alternative structures
- Not all components of a record may be used at the same time
- E.g.: Plane altitude and location on ground
- C: union
- Union is unsafe as it allows access to any member
- Pascal attempts to solve this security problem
- Access only members allowed by tag field
- Initialization not required after tag value change, so type system can be circumvented after all.


## Variant Record Example

type plane = record
flight: 0..999;
equipment: (B727, A343, B747);
case status: (inAir, taxi, atTerminal) of
inAir: (
altitude: 0..999999;
0..359);
taxi:
location: airport runway: runwayNumber);
atTerminal: (
parked: airport
gate: 1..100)
end

## Pointers

- Pascal provides typed pointers, which are more secure than untyped ones
var $p$ : $\uparrow$ real;
x : real;
c: char;
begin
new ( p ) ;
$\mathrm{p} \uparrow:=3.14159$;
$c:=p \uparrow ; \quad\{$ Illegal!
end;
- If P was untyped ( p : ¡pointer), assignment to c would be allowed (and meaningless)


## Type Equivalence

- Type checking requires that only variables with identical types can be compared/assigned to each other
- What does 'identical' mean?
- Structural equivalence
- Types having the same structure are identical
var x : record id: integer; w : real end;
var $y$ : record id: integer; w: real end;
- Name equivalence
- Types having the same name are identical


## Structural equivalence

```
type person = record id:integer; weight real; end
type car = record id:integer; weight real; end
var x: person;
var y: car;
x:= y;
    Legal by structural equivalence
    _ Probably don't want
    - Name equivalence fixes this - person and car are different names
```


## Name Structures

- Name binding mechanisms in Pascal
- Constant bindings
- Type bindings
- Variable bindings
- Procedure and function bindings
- Implicit enumeration bindings
- Label bindings


## Constants

- Pascal introduces constant declarations
const <name>=<constant>;
const MaxArray = 100;
- Allows the naming of constants in program
- Numbers should not be used in programs
- Application of Abstraction Principle


## Constants - Limitations

- Constant cannot be described by an expression
- Illegal:
const MaxArray = MaxData - 1;
- Expressions are not allowed in variable and type declarations
- Illegal:
$\operatorname{var} \mathrm{A}$ : array [0.. MaxData - 1] of real;


## Procedure Constructor

- Procedure declaration in Pascal has a strict structure
procedure <name> (<formals>)
<label declarations>
<const declarations>
<type declarations>
<var declarations>
<procedure and function declarations>
begin
end
- Similar to Algol's
- Scope essentially the same

Declarations: entire block including declarations and statements
Formals: local declarations and statements

- Names bound before they are used to support one-pass compilation

