

# CS 4100 LISP

April 22, 2011  
Based on slides by Istvan Jonyer  
Book by MacLennan  
Chapters 9, 10, 11

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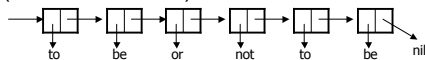
## Special Property: apval

- Assigning a value to an atom
  - (set 'Europe '(England France ...))
  - is the same as
  - (putprop 'Europe '(England France ...) 'apval)
  - 'Applied value' points to the list the atom is bound to

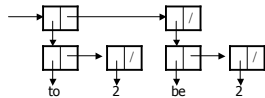
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## List Representation

- Lists are represented as linked lists  
(to be or not to be)



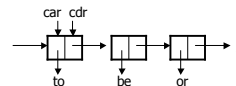
((to 2) (be 2))



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## Origins of car and cdr

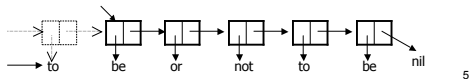
- First LISP was designed for the IBM 704
  - 1 word had 2 fields
    - Address field
    - Decrement field
  - car: “Content of Address part of Register”
  - cdr: “Content of Decrement part of Register”



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## Implementation of cons

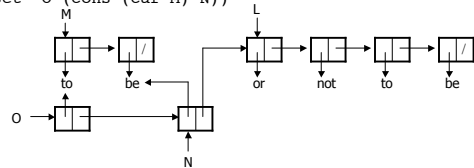
- car and cdr simply return the respective parts of the register
- cons has the job of constructing a new register using two pointers
  - Allocate new memory location
  - Fill in left and right parts of new location (cons 'to '(be or not to be))



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## Sublists Can Be Shared

```
(set 'L '(or not to be))
(set 'M '(to be))
(set 'N (cons (cadr M) L))
(set 'O (cons (car M) N))
```



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```
[10]> (set 'L '(or not to be))
(OR NOT TO BE)
[11]> (set 'M '(to be))
(TO BE)
[12]> (set 'N (cons (cadr M) L))
(BE OR NOT TO BE)
[13]> (set 'O (cons (car M) N))
(TO BE OR NOT TO BE)
```

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## List Structures Can Be Modified

- Functions discussed so far do not modify lists
- Modifying lists is possible via
  - replaca (replace address part)
  - replacd (replace decrement part)
- It is possible that more than one symbol points to a list
  - which can be modified using replaca and replacd
  - This can cause unexpected problems (like equivalence in Fortran)

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## Iteration by Recursion

- Iteration is done by recursion
  - Iteration is mostly needed to perform an operation on every element of a list
    - This can be done using combination of
      - testing for end of list,
      - operating on first element, and
      - recursing on rest of the list
- ```
(defun plus-red (a)
  (if (null a) nil
      (plus (car a) (plus-red (cdr a))) ))
```
- Notice: No array bounds are needed! Function is very general

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## Iteration = Recursion

- Theoretically, recursion and iteration have the same power, and are equivalent
- One can be translated to the other (although may not be practical)
  - Recursion → iteration
    - Use iteration and keep track of auxiliary information in an explicit stack
  - Iteration → recursion
    - Need to pass control information (variables)

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## Storage Reclamation

- What happens to *cons*'d pointers that are no longer in use?
- Explicit reclamation is the obvious / traditional way
  - C: malloc, calloc, realloc, free
  - C++: new, delete
  - Pascal: new, dispose
- Issues
  - Complicates programming
    - Requires the programmer to keep track of pointers
  - Violates security of the environment
    - Memory freed, but still referenced (dangling pointers) 11

## Automatic Storage Reclamation

- It would be nice for the system to automatically 'reclaim' storage no longer used
- System can keep track of number of references to storage
  - When references decrease to 0, storage is returned to 'free-list'
- Advantage:
  - Storage reclaimed immediately as last reference is destroyed
- Disadvantage:
  - Cyclic structures (points to itself) cannot be reclaimed

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## Garbage Collection

- A different approach is garbage collection
  - Do not keep track of references to location
  - When last reference is destroyed, we still do not do anything, and leave the memory as garbage (unused, non-reusable storage, littering the memory)
  - Collect garbage if system runs out of storage
    - Mark all areas unused
    - Then examine all visible pointers and mark storage they point to as 'used'
    - Leftover is garbage, and can be put on free-list
  - This is called the *mark-and-sweep* method

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## Garbage Collection

- Advantages
  - Fast until runs out of memory
  - No additional memory is needed for tracking references
- Disadvantages
  - Garbage collection itself can be slow
    - If memory is large, and have many references
    - Must halt entire system, since all dynamic memory must be marked as unused first
- Java uses this approach

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