## Algol Part 4

CS4100
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Based on slides by Istvan Jonyer

## Procedures are Recursive

- Recursive definitions are frequent in math and science
- Define thing in terms of itself
- Example:
- Factorial: n ! =
$\begin{cases}n^{*}(n-1)! & \text { if } n>0 \\ 1 & \text { if } n=0\end{cases}$
- Algol permits recursive procedures integer procedure fac ( n ); value $n$; integer $n$;
fac := if $n=0$ then 1 else $n$ *fac $(n-1)$;
$-\mathrm{n}=0$ ' is called the stopping condition


## Implementing Recursion

- What happens to local variable $n$ on recursive call?
- $\mathrm{fac}(3)$ is called, then $\mathrm{fac}(2)$, then $\mathrm{fac}(1)$, then fac(0)
- Would location holding 3 be overwritten?
- Yes, if same activation record was used - Solution:
- Create new activation record for each invocation of fac()


## Parameter Passing

- Modes in Algol
- Pass by value
- Pass by name
- Two modes attempt to distinguish between input only and input/output parameters


## Pass by Value

## integer procedure fac ( n );

value $n$; integer $n$;

- First part of pass by value-result (in Fortran)
- Actual copied into variable corresponding to formal
- Secure; local variable will not overwrite actual parameter
- Does not allow output parameters (input only)
- Inefficient for arrays (or other non-primitive data structures, in general)
- Copy must be made of entire array in activation record
- Copying takes time


## Pass by Name

- Based on substitution
- Consider
integer procedure Inc(n); integer n ; $\mathrm{n}:=\mathrm{n}+1$;
- And the call Inc (i)
- We need output parameter that will effect $i$, not just local $n$
- Acts like $i$ is substituted for $n$
- i := i + 1


## Copy Rule

- Procedure can be replaced by its body with actuals substituted for formals
- Revised Report 4.7.3
- Body of Inc (n)
-i := i + 1
-A[k] := A[k] + 1
- Not how it is implemented


## Pass by Name is Powerful

- Evaluate the following using pass by value, reference, and name

$$
\text { procedure } S(e l, k) \text {; }
$$

integer el, k;
begin
$\mathrm{k}:=2 ;$
el := 0;
end
$\mathrm{A}[1]:=\mathrm{A}[2]:=1$;
i $:=1$;
S(A[i], i)

- Value $A[1]=1, A[2]=1, i=1$
- Reference $A[1]=0, A[2]=1, i=2$
- Name $A[1]=1, A[2]=0, i=2$


## "Thunks"

- Implementing pass by name
- Passing the text?
- Would need to compile at runtime - not possible
- Copying compiled code?
- Would increase size of code...
- Solution: "Thunks"
- Pass address to compiled code
- Address of memory location is returned to callee to use as variable


## Pass by Name is Dangerous!

```
procedure Swap(x, y);
        integer x, y;
        begin integer t;
            t := x;
            x := y;
            y := t;
        end
```

- What is the effect of
- Swap (A[i], i)?
- Swap(i, A[i])?

- Swap(A[i], i) where A[i]=27, i=1
procedure $\operatorname{Swap}(x, y)$;
integer $\mathrm{x}, \mathrm{y}$;
begin integer $t$;
t : = A[i]; t=27
A[i] := i; A[1]=1
i $:=$ t; $\quad i=27$
end
- Swap(i, A[i]), where i=1, A[i]=27
- Swap(i, A[i]), where i=1, A[i]=27
procedure Swap(x, y);
procedure Swap(x, y);
integer x, y;
integer x, y;
begin integer t;
begin integer t;
t := i; }\quad{\begin{array}{ll}{t=1}<br>{i:=A[i];}\&{i=27}<br>{A[i] := t;}\&{A[27]=1}<br>{}
t := i; }\quad{\begin{array}{ll}{t=1}<br>{i:=A[i];}\&{i=27}<br>{A[i] := t;}\&{A[27]=1}<br>{}
end
end


## Pass-by-name

- It can be shown that there is no way to define swap in Algol-60 that works for all parameters
- Design mistake when a simple (common) procedure has such surprising properties


## Parameter Passing Modes

## Out-of-Block GOTOs

A: begin array $x[1: 100]$;

B: begin array $y[1: 100]$;
goto exit;

$$
\cdots
$$

exit:
end

- What happens to activation records?
- Program continues in different block


## Even worse...

```
    begin
        procedure P(n);
        value n; integer n;
                if n = 0 then goto out
                else P(n-1);
            P(25);
        out:
    end
- Indefinite number of activation records...
```


## Feature Interaction

- Example:
- GOTOs are simple
- Recursion is simple
- Combination is very messy
- In theory, each feature must be tested with every other one to avoid unintended consequences
- 100 features:
- Every pair: $100 \times 100=10,000$ combinations
- Every three: $100^{3}=1,000,000$ - ...
- <for statement> ::= <for clause> <statement> |
<label>: <for statement>
- <for clause> ::= for <variable> := <for list> do
- <for list> ::= <for list element> | <for list> , <for list element>
- <for list element> ::= <arithmetic expression> |
<arithmetic expression> step <arithmetic expression>
until <arithmetic expression> |
<arithmetic expression> while <Boolean expression>
- for $q:=1$ step s until $n$ do $A[q]:=B[q]$
- for $\mathrm{k}:=1, \mathrm{~V} 1 \times 2$ while $\mathrm{V} 1<\mathrm{N}$ do for $\mathrm{j}:=1+\mathrm{G}, \mathrm{L}, 1$ step 1 until $\mathrm{N}, \mathrm{C}+\mathrm{D}$ do $A[k, j]:=B[k, j]$
$\qquad$
for var $:=\exp$ step exp2 until exp3 do stat
for var := exp while exp2 do stat
- Expressions can be any arithmetic expression, including
- for $i:=i / 2$ while $i>1$ do stat
- Lists
- Conditional expressions (vs. conditional statements!) - for days := 31,
if mod(year, 4) $=0$ then 29 else 28,
$28,31,30,31,30$ do stat
- Combinations of above
- for $1:=3,7$,
$1 / 2$ while $i>1$,
do stat


## The for-loop is Very General

## Baroque Features

- Fascination-oriented features of little use
- They did it because they could
- Getting away from assembly languages as far as possible
- Baroque takes on pejorative meaning


## Baroque

- 1 : of, relating to, or having the characteristics of a style of artistic expression prevalent especially in the 17 th century that is marked generally by use of complex forms, bold ornamentation, and the juxtaposition of contrasting elements often conveying a sense of drama, movement, and tension
- 2 : characterized by grotesqueness, extravagance, complexity, or flamboyance
- 3 : irregularly shaped —used of gems <a baroque pearl>
- baroque. (2009). In Merriam-Webster Online Dictionary. Retrieved April 28, 2009, from http://www.merriamwebster.com/dictionary/baroque


## Handling Cases: switch

begin
switch wageStatus $=$ fulltime, parttime, hourly;
goto wageStatus[i];
fulltime: ...handle fulltime case...
parttime: ...handle parttime case...
parttime: goto done;
hourly: $\quad .$. handle hourly case...
goto done;
done
end;

- Elaboration on computed GOTO of Fortran (and IBM 650)
- Confusing, since switch, goto, and labels can be anywhere in the program


## Machine Independence

- Get away from formats tied to particular computers, punch cards -> free format
- How should a program be formatted?
- Left justify, one statement per line
- Like English sentence
- Structured (hierarchical) - Obeys structure principle
- Most languages followed Algol in free format


## Machine Independence

- Representation issues
- Different hardware
- Input devices
- Character sets
- Different conventions
- Math vs cs
- American vs European
- Comma (European) vs point (American) almost defeats Algol


## Compromise

- Three representations
- Reference language used in language specifications - E.g. "up arrow"
- Publication language used in publications
- E.g. sub- and super-scripts
- Hardware language to be used by implementers
- Use appropriate character set
- I/O for the computer system


## Machine Independence

- Theorem: The more trivial the point the more vehemently people will fight over it
- Which symbols
- Only those available in all sets
- Too limiting
- Independent of particular sets
- chosen


## Lexical Conventions

- Reserved words
- Cannot be used as identifiers
- Most languages
- Key words
- Words used by language are marked
- E.g. Different font or bold
- Hard to type
- Algol
- Keywords in context


## Keywords in context

- FORTRAN
- Words used by language are only keywords in context where expected
- Hard to catch errors
- Legal in PL/I

IF IF THEN
THEN $=0$;
ELSE
ELSE $=0$;

## Some Design Considerations

- From David Billington, The Tower and the Bridge 1993
- Techological Activities
- Values
- Efficiency
- Economy
- Elegance
- Dimensions
- Scientific
- Social
- Symbolic


## Efficiency

- Materials used
- Scientific Issue
- Memory
- Time
- Programmer
- Compiler
- Run


## Economy

- Cost-benefit
- Social Issue
- Benefit to programming community
- Cost: trade-offs
- Computer vs programmer time
- Increasing cost of residual bugs
- Program maintenance vs development


## Economy

- Social Influences
- Manufacturer support
- Prestigious universities teach
- Approved by influential organizations
- Standardized
- Used by "real" programmers
- Monetary values are unstable as is social climate


## Elegance

- Under-engineered
- Risk of unanticipated interactions
- Over-engineered
- Inefficient or uneconomical
- Can't always rely solely on mathematical analysis
- Always incomplete
- Simplifications
- assumptions
- Simplifications
- assumptions


## Elegance

- General Principle: Designs that look good are good
- Function follows form
- But needs to be deep (not superficial)
- Should be a joy to use
- Comfortable and safe


## Elegance

- Aesthetics comes from experience
- Design obsessively
- Criticize
- Revise
- Discard


## In Summary, Algol

- Never had widespread use
- No I/O
- Competing directly with FORTRAN
- Major milestones
- Block-structured
- Nested
- Recursive
- Free-form
- BNF - mathematical theory of formal languages


## Algol by reputation

- General
- Regular
- Elegant
- Orthogonal


## Second Generation

- Elaborations and generalizations of first generation
- Strong typing of built-in types
- Name structures hierarchically nested
- Structured control structures
- Recursion
- Parameter passing
- Syntactic structures
- Machine independent
- Moving away from fixed formats

